

# The classification of 231-avoiding permutations by descents and maximum drop

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## Abstract

We study the number of 231-avoiding permutations with  $j$ -descents and maximum drop is less than or equal to  $k$  which we denote by  $a_{n,231,j}^{(k)}$ . We show that  $a_{n,231,j}^{(k)}$  also counts the number of Dyck paths of length  $2n$  with  $n - j$  peaks and height  $\leq k + 1$ , and the number of ordered trees with  $n$  edges,  $j + 1$  internal nodes, and of height  $\leq k + 1$ . We show that the generating functions for the  $a_{n,231,j}^{(k)}$ s with  $k$  fixed satisfy a simple recursion. We also use the combinatorics of ordered trees to prove new explicit formulas for  $a_{n,231,j}^{(k)}$  as a function of  $n$  in a number of special values of  $j$  and  $k$  and prove a simple recursion for the  $a_{n,231,j}^{(k)}$ s.

**Keywords:** permutation statistics, 231-avoiding permutations, descents, drops, trees, Dyck paths.

## 1 Introduction

In [2], Chung, Claesson, Dukes, and Graham studied generating functions for permutations according to the number of descents and the maximum drop. Here if  $\sigma = \sigma_1 \dots \sigma_n$  is a permutation in the symmetric group  $S_n$ , then we say that  $\sigma$  has drop at  $i$  if  $\sigma_i < i$  and  $\sigma$  has a descent at  $i$  if  $\sigma_i > \sigma_{i+1}$ . MacMahon proved that the number of permutations with  $k$  descents is equal to the number of permutations with  $k$  drops. Let  $[n] = \{1, 2, \dots, n\}$ . We let  $\text{DES}(\sigma) = \{i \in [n] : \sigma_i > \sigma_{i+1}\}$ ,  $\text{des}(\sigma) = |\text{DES}(\sigma)|$ , and  $\text{maxdrop}(\sigma) = \max\{i - \sigma_i : i \in [n]\}$ . We let  $\mathcal{B}_n^{(k)}$  denote the set of permutations  $\sigma \in S_n$  such that  $\text{maxdrop}(\sigma) \leq k$ .

There is another interpretation of  $\mathcal{B}_n^{(k)}$  in terms of the classic bubble sort, which we denote by **bubble**. Let  $\text{bsc}(\sigma) = \min\{i : \text{bubble}^i(\sigma) = \text{id}\}$ , i.e.  $\text{bsc}(\sigma)$  is the minimum number of times that **bubble** must be applied to  $\sigma$  in order to reach the identity permutation. An inductive argument shows that  $\text{bsc}(\sigma) = \text{maxdrop}(\sigma)$ , thus  $\mathcal{B}_n^{(k)}$  is the set of permutations

in  $S_n$  which can be sorted by applying **bubble**  $k$  times. Additionally, the permutations in  $\mathcal{B}_n^{(k)}$  are in bijective correspondence with certain juggling sequences (see [2]).

Let

$$A_n^{(k)}(x) = \sum_{\sigma \in \mathcal{B}_n^{(k)}} x^{\text{des}(\sigma)} = \sum_{j=0}^{n-1} a_{n,j}^{(k)} x^j.$$

Note that for  $k \geq n-1$ ,  $\mathcal{B}_n^{(k)} = S_n$  and  $A_n^{(k)}(x)$  becomes the classic Eulerian polynomial

$$A_n(x) = \sum_{\sigma \in S_n} x^{\text{des}(\sigma)} = \sum_{j=0}^{n-1} a_{n,j} x^j.$$

The coefficient  $a_{n,j}$  is the number of permutations in  $S_n$  with  $j$  descents. These coefficients are called Eulerian numbers. For convenience we let  $A_0(x) = 1$ .

In [2], the authors show that for  $n \geq 0$ ,  $A_n^{(k)}(x)$  satisfies the following recurrence

$$A_{n+k+1}^{(k)}(x) = \sum_{i=1}^{k+1} \binom{k+1}{i} (x-1)^{i-1} A_{n+k+1-i}^{(k)}(x),$$

with the initial conditions  $A_i^{(k)}(x) = A_i(x)$  for  $0 \leq i \leq k$ . This recurrence is equivalent to the following generating function formula

$$A^{(k)}(x, t) = \sum_{n \geq 0} A_n^{(k)}(x) t^n = \frac{1 + \sum_{r=1}^k (A_r(x) - \sum_{i=1}^r \binom{k+1}{i} (x-1)^{i-1} A_{r-i}(x)) t^r}{1 - \sum_{i=1}^{k+1} \binom{k+1}{i} t^i (x-1)^{i-1}}.$$

They also find an explicit formula for  $a_{n,j}^{(k)}$ . Let

$$P_k(u) = \sum_{r=0}^k A_{k-r}(u^{k+1}) (u^{k+1} - 1)^r \sum_{i=r}^k \binom{i}{r} u^{-i},$$

and let

$$\sum_r \beta_k(r) u^r = P_k(u) \left( \frac{1 - u^{k+1}}{1 - u} \right)^{n-k},$$

then

$$A_n^{(k)}(x) = \sum_j \beta_k((k+1)j) x^j.$$

In other words, the coefficients  $a_{n,j}^{(k)}$  of the polynomial  $A_n^{(k)}(x)$  have the remarkable property that they are given by every  $(k+1)$ -st coefficient in the polynomial

$$P_k(u)(1 + u + u^2 + \cdots + u^k)^{n-k}.$$

For example setting  $n = 4$  and  $k = 2$  we have

$$P_2(u)(1 + u + u^2)^{4-2} = (1 + u + 2u^2 + u^3 + u^4)(1 + u + u^2)^2$$

$$= 1 + 3u + 7u^2 + 10u^3 + 12u^4 + 10u^5 + 7u^6 + 3u^7 + u^8.$$

So the coefficients of  $A_4^{(2)}(x)$  are given by every third coefficient in the above polynomial, that is

$$A_4^{(2)}(x) = 1 + 10x + 7x^2.$$

We now turn our attention to pattern avoidance. Given a sequence  $\sigma = \sigma_1 \dots \sigma_n$  of distinct integers, let  $\text{red}(\sigma)$  be the permutation found by replacing the  $i$ -th smallest integer that appears in  $\sigma$  by  $i$ . For example, if  $\sigma = 2754$ , then  $\text{red}(\sigma) = 1432$ . Given a permutation  $\tau = \tau_1 \dots \tau_j$  in the symmetric group  $S_j$ , we say that the pattern  $\tau$  *occurs* in  $\sigma = \sigma_1 \dots \sigma_n \in S_n$  provided there exists  $1 \leq i_1 < \dots < i_j \leq n$  such that  $\text{red}(\sigma_{i_1} \dots \sigma_{i_j}) = \tau$ . We say that a permutation  $\sigma$  *avoids* the pattern  $\tau$  if  $\tau$  does not occur in  $\sigma$ . Let  $S_n(\tau)$  denote the set of permutations in  $S_n$  which avoid  $\tau$ . In the theory of permutation patterns (see [6] for a comprehensive introduction to the area),  $\tau$  is called a *classical pattern*. We let  $\mathcal{B}_{n,\tau}^{(k)} = S_n(\tau) \cap \mathcal{B}_n^{(k)}$ . Thus  $\mathcal{B}_{n,\tau}^{(k)}$  is the set of  $\sigma \in S_n$  such that  $\text{maxdrop}(\sigma) \leq k$  and  $\sigma$  avoids  $\tau$ . For  $k \geq 1$ , we let  $\mathcal{E}_{n,\tau}^{(k)} = \mathcal{B}_{n,\tau}^{(k)} - \mathcal{B}_{n,\tau}^{(k-1)}$ . Thus  $\mathcal{E}_{n,\tau}^{(k)}$  is the set of  $\sigma \in S_n$  such that  $\text{maxdrop}(\sigma) = k$  and  $\sigma$  avoids  $\tau$ . We let

$$\begin{aligned} A_{n,\tau}^{(k)}(x) &= \sum_{\sigma \in \mathcal{B}_{n,\tau}^{(k)}} x^{\text{des}(\sigma)} = \sum_{j=0}^{n-1} a_{n,\tau,j}^{(k)} x^j \text{ and} \\ E_{n,\tau}^{(k)}(x) &= \sum_{\sigma \in \mathcal{E}_{n,\tau}^{(k)}} x^{\text{des}(\sigma)} = \sum_{j=0}^{n-1} e_{n,\tau,j}^{(k)} x^j. \end{aligned}$$

Let

$$A_\tau^{(k)}(x, t) = 1 + \sum_{n \geq 1} A_{n,\tau}^{(k)}(x) t^n \quad (1)$$

and

$$E_\tau^{(k)}(x, t) = 1 + \sum_{n \geq 1} E_{n,\tau}^{(k)}(x) t^n \quad (2)$$

Note that for  $k \geq 1$ ,  $E_{n,\tau}^{(k)}(x) = A_{n,\tau}^{(k)}(x) - A_{n,\tau}^{(k-1)}(x)$  so that  $E_{n,\tau}^{(k)}(x, t) = A_\tau^{(k)}(x, t) - A_\tau^{(k-1)}(x, t)$ .

The main goal of this paper is to study the generating functions  $A_{231}^{(k)}(x, t)$  and  $E_{231}^{(k)}(x, t)$ .

We remark that the set  $\mathcal{B}_{n,231}^{(k)}$  can also be interpreted in terms of sorting algorithms. The 231-avoiding permutations are precisely the permutations which can be sorted by one application of the stack sort, which we denote by **stack** (see [10]). So  $\mathcal{B}_{n,231}^{(k)} = \{\sigma \in S_n : \text{bubble}^k(\sigma) = \text{id}, \text{ and } \text{stack}(\sigma) = \text{id}\}$ , i.e. the permutations in  $S_n$  which can be sorted by one stack sort, but require  $k$  bubble sorts to be sorted.

Note that the only permutation  $\sigma \in S_n$  such that  $\text{maxdrop}(\sigma) = 0$  is the identity permutation  $\sigma = 123 \dots n$  which is 231-avoiding. Thus  $A_{n,231}^{(0)}(x) = 1$  for all  $n \geq 1$  so that

$$A_{231}^{(0)}(x, t) = \frac{1}{1-t}. \quad (3)$$

Our key theorem is to show that generating functions  $A_{231}^{(k)}(x, t)$  for  $k \geq 1$  satisfy the following simple recursion.

**Theorem 1.** *For all  $k \geq 1$ ,*

$$A_{231}^{(k)}(x, t) = \frac{1}{1 - t + tx - txA_{231}^{(k-1)}(x, t)} \quad (4)$$

where

$$A_{231}^{(0)}(x, t) = \frac{1}{1 - t}.$$

Theorem 1 allowed us to explicitly compute the values of  $a_{n,231,j}^{(k)}$  and  $e_{n,231,j}^{(k)}$  for small values of  $j$ ,  $k$ , and  $n$  which lead us to conjecture a number of simple formulas for  $a_{n,231,j}^{(k)}$  and  $e_{n,231,j}^{(k)}$  in certain special cases. For example, we shall show that for all  $n, j \geq 1$  and all  $k \geq j$ ,  $a_{n,231,j}^{(k)} = N(n, n-j) = \frac{1}{n} \binom{n}{j} \binom{n}{j+1}$  and  $e_{n,231,j}^{(j)} = \binom{n+j-1}{2j}$ . Here the  $N(n, j)$ s are the Narayana numbers which count the number of Dyck paths of length  $2n$  with  $j$  peaks and the number of ordered trees  $n$  edges and  $k$  leaves. This suggested that the numbers  $a_{n,231,j}^{(k)}$  should also have natural combinatorial interpretations in terms of Dyck paths and ordered trees. In fact, we construct bijections to show that  $a_{n,231,j}^{(k)}$  is the number of ordered trees with height less than or equal to  $k+1$ ,  $n$  edges, and  $j+1$  internal nodes and is the number of Dyck paths of length  $2n$  with  $n-j$  peaks and height less than or equal to  $k+1$ .

Kemp [5] gave a general formula for the number of ordered trees with height less than or equal to  $k$ ,  $n$  edges, and  $j$  internal nodes so that we have a general formula for  $a_{n,231,j}^{(k)}$ . However, in many cases, Kemp's formula is unnecessarily complicated so that we use the combinatorics of ordered trees to derive an number of elegant formulas and recursions for the  $a_{n,231,j}^{(k)}$ s. For example, we shall show that

$$a_{n,231,j}^{(j-2)} = \frac{1}{n} \binom{n}{j} \binom{n}{j+1} - \binom{n+j-1}{2j} - (2j-3) \binom{n+j-2}{2j}$$

and that the  $a_{n,231,j}^{(k)}$ s satisfy the following simple recursion:

$$a_{n,231,j}^{(k)} = \sum_{i=0}^j a_{j,231,i}^{(k-1)} \binom{n+i}{2j}.$$

The outline of this paper is as follows. In section 2, we shall prove theorem 1 as well as provide simple proofs of the fact that  $a_{n,231,1}^{(k)} = \binom{n}{2}$  for all  $k \geq 1$  and  $n \geq 2$ ,  $a_{n,231,2}^{(k)} = \frac{(n-1)^2((n-1)^2-1)}{12}$  for all  $k \geq 2$  and  $n \geq 3$ , and that  $e_{n,231,2}^{(2)} = \binom{n+1}{4}$  for all  $n \geq 3$ . In section 3, we shall prove our alternative combinatorial interpretations of  $a_{n,231,j}^{(k)}$  in terms of ordered trees and Dyck paths. In section 4, we shall use the combinatorics of ordered trees to prove a number of formulas for  $a_{n,231,j}^{(k)}$  and  $e_{n,231,j}^{(k)}$  as well as derive a new recursion for the  $a_{n,231,j}^{(k)}$ s. Finally, in section 5, we shall briefly discuss some combinatorial identities that arise by comparing Kemp's formula and our formulas.

## 2 Proof of Theorem 1

In this section, we shall prove Theorem 1. The proof proceeds by classifying the 231-avoiding permutations  $\sigma = \sigma_1 \dots \sigma_n$  by the position of  $n$  in  $\sigma$ . Clearly each  $\sigma \in S_n(231)$  has the structure pictured in Figure 1. That is, in the graph of  $\sigma$ , the elements to the left of  $n$ ,  $C_i(\sigma)$ , have the structure of a 231-avoiding permutation, the elements to the right of  $n$ ,  $D_i(\sigma)$ , have the structure of a 231-avoiding permutation, and all the elements in  $C_i(\sigma)$  lie below all the elements in  $D_i(\sigma)$ . Note that the number of 231-avoiding permutations in  $S_n$  is the Catalan number  $C_n = \frac{1}{n+1} \binom{2n}{n}$  and the generating function for the  $C_n$ 's is given by

$$C(t) = \sum_{n \geq 0} C_n t^n = \frac{1 - \sqrt{1 - 4t}}{2t} = \frac{2}{1 + \sqrt{1 - 4t}}. \quad (5)$$

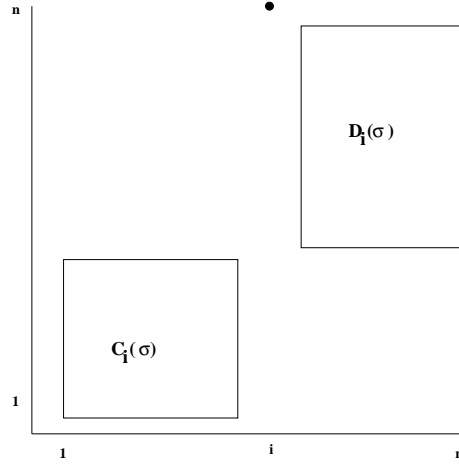


Figure 1: The structure of 231-avoiding permutations.

Suppose that  $k \geq 1$  and  $n \geq 2$ . Let  $\mathcal{B}_{n,i,231}^{(k)}$  denote the set of  $\sigma \in \mathcal{B}_{n,231}^{(k)}$  such that  $\sigma_i = n$ . Clearly if  $\sigma \in \mathcal{B}_{n,i,231}^{(k)}$ , then  $C_i(\sigma)$  must be a permutation in  $S_{i-1}(231)$  such that  $\text{maxdrop}(C_i(\sigma)) \leq k$ . Similarly,  $\tau = \text{red}(D_i(\sigma))$  must be a permutation in  $S_{n-i}(231)$  such that  $\text{maxdrop}(\tau) \leq k - 1$ . That is, we can consider  $D_i(\sigma)$  as a map from  $\{(i+1, \dots, i+(n-i))\}$  into  $\{(i-1+1, \dots, (i-1)+(n-i))\}$ . Thus for  $j = 1, \dots, n-i$ , a drop  $j - \tau_j$  in  $\tau$  corresponds to a drop  $i+j - (i-1+\tau_j) = i+j - \sigma_{i+j} - 1$  in  $\sigma$ . Thus the drop at position  $j$  in  $\tau$  is one less than the drop at position  $i+j$  in  $\sigma$ . Now if  $i \leq n-1$ , then  $\sigma_i = n$  will start a descent in  $\sigma$ . Thus the possible choices for  $C_i(\sigma)$  will contribute a factor of  $A_{i-1,231}^{(k)}(x)$  to  $\sum_{\sigma \in \mathcal{B}_{n,i,231}^{(k)}} x^{\text{des}(\sigma)}$  and the possible choices for  $D_i(\sigma)$  will contribute a factor of  $A_{n-i,231}^{(k-1)}(x)$  to  $\sum_{\sigma \in \mathcal{B}_{n,i,231}^{(k)}} x^{\text{des}(\sigma)}$ . Thus the contribution of the permutations in  $\mathcal{B}_{n,i,231}^{(k)}$  to  $A_{n,231}^{(k)}(x)$  is  $x A_{i-1,231}^{(k)}(x) A_{n-i,231}^{(k-1)}(x)$ . Finally, it is easy to see that the contribution of the permutations

in  $\mathcal{B}_{n,n,231}^{(k)}$  to  $A_{n,231}^{(k)}(x)$  is just  $A_{n-1,231}^{(k)}(x)$ . It follows that for  $n \geq 2$ ,

$$A_{n,231}^{(k)}(x) = A_{n-1,231}^{(k)}(x) + x \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x) A_{n-i,231}^{(k-1)}(x). \quad (6)$$

Note that  $A_{1,231}^{(k)}(x) = 1$  so that if we define  $A_{0,231}^{(k)}(x) = 1$ , then (6) also holds for  $n = 1$ . Multiplying both sides of (6) by  $t^n$  and summing for  $n \geq 1$ , we see that for  $k \geq 1$ ,

$$\begin{aligned} A_{231}^{(k)}(x, t) - 1 &= t A_{231}^{(k)}(x, t) + t \sum_{n \geq 1} t^{n-1} x \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x) A_{n-i,231}^{(k-1)}(x) \\ &= t A_{231}^{(k)}(x, t) + t x A_{231}^{(k)}(x, t) (A_{231}^{(k-1)}(x, t) - 1). \end{aligned}$$

Solving this equation for  $A_{231}^{(k)}(x, t)$ , we see that

$$A_{231}^{(k)}(x, t) = \frac{1}{1 - t + xt - xt A_{231}^{(k-1)}(x, t)} \quad (7)$$

which proves Theorem 1.

One can use Mathematica to calculate the first few of the generating functions  $A_{231}^{(k)}(x, t)$ .

$$\begin{aligned} A_{231}^{(0)}(x, t) &= \frac{1}{1 - t}, \\ A_{231}^{(1)}(x, t) &= \frac{1 - t}{1 - 2t + (1 - x)t^2}, \\ A_{231}^{(2)}(x, t) &= \frac{1 - 2t + (1 - x)t^2}{1 - 3t + (3 - 2x)t^2 - (1 - x)^2 t^3}, \\ A_{231}^{(3)}(x, t) &= \frac{1 - 3t + (3 - 2x)t^2 - (1 - x)^2 t^3}{1 - 4t + 3(x - 2)t^2 - 2(2 - 3x_x^2)t^3 + (1 - x)^3 t^4}, \\ A_{231}^{(4)}(x, t) &= \frac{1 - 4t + 3(x - 2)t^2 - 2(2 - 3x_x^2)t^3 + (1 - x)^3 t^4}{1 - 5t + (10 - 4x)t^2 - (10 - 12x + 3x^2)t^3 - (1 - x)^2(2x - 5)t^4 - (1 - x)^4 t^5}. \end{aligned}$$

One can also use Mathematica to find the initial terms of the generating function  $A_{231}^{(k)}(t, x)$ . For example, we have computed that

$$\begin{aligned} A_{231}^{(1)}(t, x) &= 1 + t + (1 + x)t^2 + (1 + 3x)t^3 + (1 + 6x + x^2)t^4 + (1 + 10x + 5x^2)t^5 + \\ &\quad (1 + 15x + 15x^2 + x^3)t^6 + (1 + 21x + 35x^2 + 7x^3)t^7 + \\ &\quad (1 + 28x + 70x^2 + 28x^3 + x^4)t^8 + (1 + 36x + 126x^2 + 84x^3 + 9x^4)t^9 + \\ &\quad (1 + 45x + 210x^2 + 210x^3 + 45x^4 + x^5)t^{10} + \\ &\quad (1 + 55x + 330x^2 + 462x^3 + 165x^4 + 11x^5)t^{11} + \dots \end{aligned}$$

$$\begin{aligned}
& A_{231}^{(2)}(t, x) \\
& = 1 + t + (1 + x)t^2 + (1 + 3x + x^2)t^3 + (1 + 6x + 6x^2)t^4 + (1 + 10x + 20x^2 + 3x^3)t^5 + \\
& (1 + 15x + 50x^2 + 22x^3 + x^4)t^6 + (1 + 21x + 105x^2 + 91x^3 + 15x^4)t^7 + \\
& (1 + 28x + 196x^2 + 280x^3 + 100x^4 + 5x^5)t^8 + \\
& (1 + 36x + 336x^2 + 714x^3 + 444x^4 + 65x^5 + x^6)t^9 + \\
& (1 + 45x + 540x^2 + 1596x^3 + 1530x^4 + 441x^5 + 28x^6)t^{10} + \\
& (1 + 55x + 825x^2 + 3234x^3 + 4422x^4 + 2101x^5 + 301x^6 + 7x^7)t^{11} + \dots
\end{aligned}$$

$$\begin{aligned}
& A_{231}^{(3)}(t, x) \\
& = 1 + t + (1 + x)t^2 + (1 + 3x + x^2)t^3 + (1 + 6x + 6x^2 + x^3)t^4 + \\
& (1 + 10x + 20x^2 + 10x^3)t^5 + (1 + 15x + 50x^2 + 50x^3 + 6x^4)t^6 + \\
& (1 + 21x + 105x^2 + 175x^3 + 60x^4 + 3x^5)t^7 + \\
& (1 + 28x + 196x^2 + 490x^3 + 325x^4 + 53x^5 + x^6)t^8 + \\
& (1 + 36x + 336x^2 + 1176x^3 + 1269x^4 + 428x^5 + 35x^6)t^9 + \\
& (1 + 45x + 540x^2 + 2520x^3 + 4005x^4 + 2289x^5 + 427x^6 + 15x^7)t^{10} + \\
& (1 + 55x + 825x^2 + 4950x^3 + 10857x^4 + 9394x^5 + 3122x^6 + 316x^7 + 5x^8)t^{11} + \dots
\end{aligned}$$

$$\begin{aligned}
& A_{231}^{(4)}(t, x) = \\
& 1 + t + (1 + x)t^2 + (1 + 3x + x^2)t^3 + (1 + 6x + 6x^2 + x^3)t^4 + \\
& (1 + 10x + 20x^2 + 10x^3 + x^4)t^5 + (1 + 15x + 50x^2 + 50x^3 + 15x^4)t^6 + \\
& (1 + 21x + 105x^2 + 175x^3 + 105x^4 + 10x^5)t^7 + \\
& (1 + 28x + 196x^2 + 490x^3 + 490x^4 + 130x^5 + 6x^6)t^8 + \\
& (1 + 36x + 336x^2 + 1176x^3 + 1764x^4 + 890x^5 + 128x^6 + 3x^7)t^9 + \\
& (1 + 45x + 540x^2 + 2520x^3 + 5292x^4 + 4291x^5 + 1246x^6 + 105x^7 + x^8)t^{10} + \\
& (1 + 55x + 825x^2 + 4950x^3 + 13860x^4 + 16401x^5 + 7945x^6 + 1435x^7 + 70x^8)t^{11} +
\end{aligned}$$

The generating function for the number of 231-avoiding permutations  $\sigma$  with  $\text{maxdrop}(\sigma) \leq k$  is  $A_{231}^{(k)}(1, t)$  for any  $k \geq 0$ . These are easily computed using Theorem 1 and Mathematica.

For example, we have computed that

$$\begin{aligned}
A_{231}^{(0)}(1, t) &= \frac{1}{1-t}, \\
A_{231}^{(1)}(1, t) &= \frac{1-t}{1-2t}, \\
A_{231}^{(2)}(1, t) &= \frac{1-2t}{1-3t+t^2}, \\
A_{231}^{(3)}(1, t) &= \frac{1-3t+t^2}{1-4t+3t^2}, \\
A_{231}^{(4)}(1, t) &= \frac{1-4t+3t^2}{1-5t+6t^2-t^3}, \text{ and} \\
A_{231}^{(5)}(1, t) &= \frac{1-5t+6t^2-t^3}{1-6t+10t^2-4t^3}.
\end{aligned}$$

These generating functions have recently turned up in a completely different context. In [8], Kitaev, Remmel, and Tiefenbruck studied what they called *quadrant marked mesh patterns*. That is, let  $\sigma = \sigma_1 \dots \sigma_n$  be a permutation written in one-line notation. Then we will consider the graph of  $\sigma$ ,  $G(\sigma)$ , to be the set of points  $(i, \sigma_i)$  for  $i = 1, \dots, n$ . For example, the graph of the permutation  $\sigma = 471569283$  is pictured in Figure 2. Then if we draw a coordinate system centered at a point  $(i, \sigma_i)$ , we will be interested in the points that lie in the four quadrants I, II, III, and IV of that coordinate system as pictured in Figure 2. For any  $a, b, c, d \in \mathbb{N}$  where  $\mathbb{N} = \{0, 1, 2, \dots\}$  is the set of natural numbers and any  $\sigma = \sigma_1 \dots \sigma_n \in S_n$ , we say that  $\sigma_i$  matches the quadrant marked mesh pattern  $MMP(a, b, c, d)$  in  $\sigma$  if in  $G(\sigma)$  relative to the coordinate system which has the point  $(i, \sigma_i)$  as its origin, there are  $\geq a$  points in quadrant I,  $\geq b$  points in quadrant II,  $\geq c$  points in quadrant III, and  $\geq d$  points in quadrant IV. For example, if  $\sigma = 471569283$ , the point  $\sigma_4 = 5$  matches the simple marked mesh pattern  $MMP(2, 1, 2, 1)$  since relative to the coordinate system with origin  $(4, 5)$ , there are 3 points in  $G(\sigma)$  in quadrant I, 1 point in  $G(\sigma)$  in quadrant II, 2 points in  $G(\sigma)$  in quadrant III, and 2 points in  $G(\sigma)$  in quadrant IV. Note that if a coordinate in  $MMP(a, b, c, d)$  is 0, then there is no condition imposed on the points in the corresponding quadrant. In [8], the authors studied the generating functions

$$Q_{132}^{(a,b,c,d)}(t, x) = 1 + \sum_{n \geq 1} t^n Q_{n,132}^{(a,b,c,d)}(x) \quad (8)$$

where for any  $a, b, c, d \in \mathbb{N}$ ,

$$Q_{n,132}^{(a,b,c,d)}(x) = \sum_{\sigma \in S_n(132)} x^{\text{mmp}^{(a,b,c,d)}(\sigma)}. \quad (9)$$

It turns out that  $Q^{(k,0,0,0)}(t, 0) = A^{(k-1)}(1, t)$  for all  $k \geq 2$  since it was shown in [8] that

$$Q_{132}^{(1,0,0,0)}(t, 0) = \frac{1}{1-t}$$



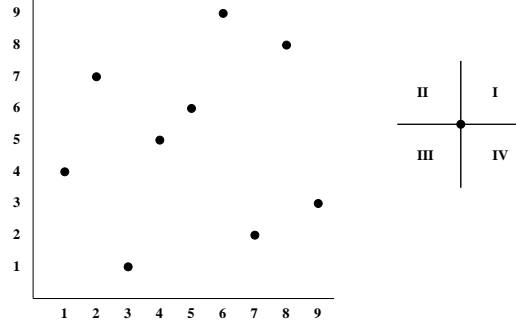


Figure 2: The graph of  $\sigma = 471569283$ .

and for  $k > 1$ ,

$$Q_{132}^{(k,0,0)}(t, 0) = \frac{1}{1 - tQ_{132}^{(k-1,0,0)}(t, 0)}.$$

Thus the number of 231-avoiding permutations with  $\text{maxdrop}(\sigma) \leq k - 1$  is the number of 132-avoiding permutations which have no occurrences of the quadrant mesh pattern  $MMP(k, 0, 0, 0)$ . In fact, one can use the recursions satisfied by  $A_{n,231}^{(k-1)}(1)$  and  $Q_{n,132}^{(k,0,0)}$  to give a bijective proof of this fact. It was also shown in [8] that the number of permutations  $\sigma \in S_n(132)$  which have no occurrences of the quadrant mesh pattern  $MMP(k, 0, 0, 0)$  is also equal to the number of Dyck paths of length  $2n$  such that all steps have height  $\leq k$ .

Many of the sequences  $(A_{n,231}^{(k)}(1))_{n \geq 1}$  as well as many of the sequences  $(a_{n,231,r}^{(k)})_{n \geq 1}$  appear in the OEIS. For example, the sequence  $(A_{n,231}^{(2)}(1))_{n \geq 0}$  starts out with 1, 1, 2, 5, 13, 34, 89, 233, 610, 1597, . . . . This is sequence A001519 in the OEIS. It immediately follows from the generating function

$$A_{231}^{(2)}(1, t) = \frac{1 - 2t}{1 - 3t + t^2}$$

that the numbers  $A_{n,231}^{(2)}(1)$  satisfy the simple recursion that

$$A_{n,231}^{(2)}(1) = 3A_{n-1,231}^{(2)}(1) - A_{n-2,231}^{(2)}(1)$$

with initial conditions that  $A_{0,231}^{(2)}(1) = A_{1,231}^{(2)}(1) = 1$ . The OEIS lists many combinatorial interpretations of these numbers including the number of permutations of  $S_{n+1}$  which avoid 321 and 3412 and the number of ordered trees with  $n + 1$  edges and height of at most 3. In section 3, we shall establish a direction connection between ordered trees and the permutations in  $\mathcal{B}_{n,231}^{(k)}$  which will explain this and many other formulas.

The sequence  $(A_{n,231}^{(3)}(1))_{n \geq 0}$  starts out with 1, 1, 2, 5, 14, 41, 122, 365, 1094, 3281, . . . . This is sequence A124302 in the OEIS. This sequence also has many combinatorial definitions including the number of set partitions of  $[n] = \{1, \dots, n\}$  of length  $\leq 3$ . The sequence  $(A_{n,231}^{(4)}(1))_{n \geq 0}$  starts out with 1, 1, 2, 5, 14, 42, 131, 417, 1341, 4434, . . . . which is sequence A080937 in the OEIS. The sequence  $(A_{n,231}^{(5)}(1))_{n \geq 0}$  starts out with 1, 1, 2, 5, 14, 42, 132, 428, 1416, 4744, . . . which is sequence A024175 in the OEIS.

We end this section, with a few simple results which can be easily proved from (6). These results will be generalized in the subsequent section when we consider a bijection between ordered trees and permutations  $\sigma$  with  $\text{maxdrop}(\sigma) \leq k$ .

**Theorem 2.** 1. For all  $n \geq 1$ ,  $A_{n,231}^{(1)}(1) = 2^{n-1}$ .

2. For all  $r \geq 1$  and  $n \geq 2r$ ,  $A_{n,231}^{(1)}(x)|_{x^r} = a_{n,231,r}^{(1)} = \binom{n}{2r}$ .

*Proof.* Part (1) follows immediately from the fact that  $A_{231}^{(1)}(1, t) = \frac{1-t}{1-2t}$ . It is also easy to give a direct inductive proof of part (1). That is, for part (1), clearly the statement holds for  $n = 1$  since  $A_{1,231}^{(1)}(1) = 1$ . But then using the fact that  $A_{n,231}^{(0)}(1) = 1$  for all  $n \geq 0$ , we see that (6) implies

$$\begin{aligned} A_{n,231}^{(1)}(1) &= A_{n-1,231}^{(1)}(1) + \sum_{i=1}^{n-1} A_{i-1,231}^{(1)}(1) A_{n-i,231}^{(0)}(1) \\ &= A_{n-1,231}^{(1)}(1) + A_{n-1,231}^{(0)}(1) + \sum_{i=2}^{n-1} A_{i-1,231}^{(1)}(1) \\ &= 2^{n-2} + 1 + \sum_{i=2}^{n-1} 2^{i-2} = 2^{n-1}. \end{aligned}$$

In fact, we can directly construct all the elements  $\mathcal{B}_{n,231}^{(1)}$ . We let  $[n] = \{1, \dots, n\}$  and, if  $1 \leq i < j \leq n$ , we let  $[i, j] = \{s \in [n] : i \leq s \leq j\}$  be the interval from  $i$  to  $j$ . Let  $\mathcal{P}([n])$  denote the set of all subsets of  $[n]$  and  $\mathcal{P}_e([n])$  denote the set of all elements of  $\mathcal{P}([n])$  that have even cardinality. Clearly, the cardinality of  $\mathcal{P}_e([n])$  is  $2^{n-1}$ . We define bijection  $\phi : \mathcal{P}_e([n]) \rightarrow \mathcal{B}_{n,231}^{(1)}$ . We let  $\phi(\emptyset) = 12 \dots n$ . Now if  $S = \{s_1, s_2, \dots, s_{2r-1}, s_{2r}\} \in \mathcal{P}_e([n])$  where  $1 \leq s_1 < s_2 < \dots < s_{2r-1} < s_{2r} \leq n$ , then we consider the intervals  $I_j = [s_{2j-1}, s_{2j}]$  for  $j = 1, \dots, r$ . We define  $\phi(S) = \tau^S = \tau_1^S \dots \tau_n^S$  to be the permutation in  $S_n$  such that  $\tau_i^S = i$  if  $i$  is not in one of the intervals  $I_1, \dots, I_r$ , and  $\tau_{s_{2j-1}}^S \dots \tau_{s_{2j}}^S = s_{2j} s_{2j-1} (s_{2j-1} + 1) \dots (s_{2j} - 1)$ . For example, if  $n = 12$ , and  $S = \{1, 3, 6, 8, 10, 12\}$ , then  $I_1 = [1, 3]$ ,  $I_2 = [6, 8]$ , and  $I_3 = [10, 12]$ . Thus

$$\tau^S = 3 \ 1 \ 2 \ 4 \ 5 \ 8 \ 6 \ 7 \ 9 \ 12 \ 10 \ 11.$$

Note on each of the intervals  $I_j$ ,  $\tau^S$  has maximum drop 1 so that  $\text{maxdrop}(\tau^S) \leq 1$  for all  $S \in \mathcal{P}_e([n])$ . Moreover it is easy to see that  $\tau^S$  is 231-avoiding and that we can recover  $S$  from  $\tau^S$ . Thus  $\phi$  is a one-to-one map from  $\mathcal{P}_e([n])$  into  $\mathcal{B}_{n,231}^{(1)}$ . However, since we know that  $|\mathcal{P}_e([n])| = |\mathcal{B}_{n,231}^{(1)}|$ ,  $\phi$  must also be a surjection. Thus  $\phi$  is a bijection from  $\mathcal{P}_e([n])$  onto  $\mathcal{B}_{n,231}^{(1)}$ .

Note that  $\phi$  also has the property that for all  $S \in \mathcal{P}_e([n])$ ,  $\text{des}(\phi(S)) = \frac{|S|}{2}$ . Thus it follows that the number of  $\sigma \in \mathcal{B}_{n,231}^{(1)}$  such that  $\text{des}(\sigma) = r$  equals the number of subsets  $S$  of  $[n]$  of size  $2r$ . That is,  $A_{n,231}^{(1)}(x)|_{x^r} = \binom{n}{2r}$  which proves part (2).  $\square$

In fact, our construction in Theorem 2 constructs all the possible elements of  $S_n(231)$  with exactly one descent. That is, we claim that if  $\sigma \in S_n(231)$  and  $\text{des}(\sigma) = 1$ , then  $\text{maxdrop}(\sigma) = 1$ . Suppose that  $\sigma = \sigma_1 \dots \sigma_n \in S_n(231)$ ,  $\text{des}(\sigma) = 1$  and  $\text{maxdrop}(\sigma) = k$  where  $k \geq 2$ . Let  $i$  be least element such that  $i - \sigma_i = k$ . Thus  $\sigma_i = i - k$ . Our choice of  $i$  ensures that  $\sigma_{i-1} \geq i - k$  since we can have a drop of at most  $k - 1$  at position  $i - 1$ . But since  $\sigma_i = i - k$ , it must be the case that  $\sigma_{i-1} > i - k = \sigma_i$ . Thus the only descent of  $\sigma$  must occur at position  $i - 1$ . This means that  $\sigma_1 < \dots < \sigma_{i-1}$ . Since  $k \geq 2$ , there must be at least two elements in  $\sigma_1 \dots \sigma_{i-1}$  which are greater than or equal to  $i - k$  which would mean that there is an occurrence of 231 in  $\sigma_1 \dots \sigma_i$ . Thus if  $\sigma \in S_n(231)$  and  $\text{des}(\sigma) = 1$ , then it must be the case that  $\text{maxdrop}(\sigma) = 1$ . Thus the elements of the form  $\phi(S)$  where  $|S| = 2$  consists of all the elements of  $S_n(231)$  such that  $\text{des}(\sigma) = 1$ . It thus follows that we have the following theorem.

**Theorem 3.** *For all  $k \geq 1$  and  $n \geq 2$ ,*

$$A_{n,231}^{(k)}(x)|_x = a_{n,231,1}^{(k)} = \binom{n}{2}.$$

The sequence  $(a_{n,231,2}^{(2)})_{n \geq 3}$  starts out with 1, 6, 20, 50, 105, 196, 336, 540, .... This is sequence A002415 in the OEIS whose  $n$ -th term is  $\frac{n^2(n^2-1)}{12}$ . These numbers are known as the 4-dimensional pyramidal numbers. They have several combinatorial interpretations including the number of squares with corners in the  $n \times n$  grid. We can prove the following general theorem about such numbers.

**Theorem 4.** *For all  $k \geq 2$  and  $n \geq 3$ ,  $A_{n,231}^{(k)}(x)|_{x^2} = \frac{(n-1)^2((n-1)^2-1)}{12}$ .*

*Proof.* We proceed by induction on  $n$ . For the base case, note that the only permutation  $\sigma \in S_n(231)$  with 2 descents is 321 which has maximum drop 2. It follows that  $A_{3,231}^{(k)}(x)|_{x^2} = 1$  for all  $k \geq 2$  so that our formulas hold for  $n = 3$ .

Next assume that  $n > 3$  and our formula holds for all  $m < n$ . Then by (6),

$$A_{n,231}^{(k)}(x)|_{x^2} = A_{n-1,231}^{(k)}(x)|_{x^2} + \left( \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x) A_{n-i,231}^{(k-1)}(x) \right) |_x. \quad (10)$$

By induction,  $A_{n-1,231}^{(k)}(x)|_{x^2} = \frac{(n-2)^2((n-2)^2-1)}{12}$ . Note that

$$\begin{aligned} \left( \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x) A_{n-i,231}^{(k-1)}(x) \right) |_x &= \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x)|_x A_{n-i,231}^{(k-1)}(x)|_{x^0} + \\ &\quad \sum_{i=1}^{n-1} A_{i-1,231}^{(k)}(x)|_{x^0} A_{n-i,231}^{(k-1)}(x)|_x. \end{aligned}$$

But we know that  $A_{i-1,231}^{(k)}(x)|_{x^0} = 1$  for all  $k$  and  $n$  since the only permutation with no descents is the identity. By Theorem 3, we know that  $A_{n,231}^{(k)}(x)|_x = \binom{n}{2}$  for all  $k \geq 1$  and

$n \geq 1$ . It follows that

$$\begin{aligned}
A_{n,231}^{(k)}(x)|_{x^2} &= \frac{(n-2)^2((n-2)^2-1)}{12} + \sum_{i=1}^{n-1} \binom{i-1}{2} + \sum_{i=1}^{n-1} \binom{n-i}{2} \\
&= \frac{(n-2)^2((n-2)^2-1)}{12} + \binom{n-1}{3} + \binom{n}{3} \\
&= \frac{(n-1)^2((n-1)^2-1)}{12}
\end{aligned}$$

where the last equality can be checked in Mathematica.  $\square$

Theorems 2 and Theorem 4 automatically imply the following theorem.

**Theorem 5.** For  $n \geq 3$ ,  $E_{n,132}^{(2)}(x)|_{x^2} = \binom{n+1}{4}$ .

*Proof.* By Theorem 2, we know that  $A_{n,132}^{(1)}(x)|_{x^2} = \binom{n}{4}$  and Theorem 4, we know that  $A_{n,132}^{(2)}(x)|_{x^2} = \frac{(n-1)^2(((n-1)^2)-1)}{12}$ . Thus

$$\begin{aligned}
E_{n,132}^{(2)}(x)|_{x^2} &= A_{n,132}^{(2)}(x)|_{x^2} - A_{n,132}^{(1)}(x)|_{x^2} \\
&= \frac{(n-1)^2(((n-1)^2)-1)}{12} - \binom{n}{4} = \binom{n+1}{4}.
\end{aligned}$$

$\square$

### 3 Ordered trees of bounded height

In this section we show there is a bijective correspondence between permutations in  $S_n(231)$  with a given maximum drop and a given number of descents, to a certain class of trees. An ordered tree is a rooted tree where the children of each vertex are ordered, so for example we can refer to the left-most child of a vertex. We use the convention of placing the root at the top of the tree. Micheli and Rossin show there is a bijection between 231-avoiding permutations and ordered trees [12]. Here we show this same bijection also carries additional information about the descents and maximum drop size of 231-avoiding permutations. The level of a vertex is the distance of the shortest path from that vertex to the root. The height of an ordered tree is the maximum of the levels of all vertices in the tree. An internal node is a vertex which has at least one child. Let  $\mathcal{T}_{n,j}^{(k)}$  denote the set of all ordered trees having  $n$  edges, height less than or equal to  $k$ , and  $j$  internal nodes. Let  $\mathcal{B}_{n,231,j}^{(k)}$  denote the set of permutations in  $\sigma \in S_n(231)$  with  $\text{des}(\sigma) = j$  and  $\text{maxdrop}(\sigma) \leq k$ , thus  $|\mathcal{B}_{n,231,j}^{(k)}| = A_{n,231}^{(k)}(x)|_{x^j} = a_{n,231,j}^{(k)}$  (not to be confused with  $\mathcal{B}_{n,i,231}^{(k)}$  in the proof of Theorem 1).

**Theorem 6.** There is a bijection  $\phi : \mathcal{T}_{n,j}^{(k)} \rightarrow \mathcal{B}_{n,231,j-1}^{(k-1)}$ , for all  $n, k, j \geq 1$ , thus

$$a_{n,231,j}^{(k)} = |\mathcal{T}_{n,j+1}^{(k+1)}|.$$

In other words  $a_{n,231,j}^{(k)}$  is equal to the number of ordered trees with  $n$  edges,  $j+1$  internal nodes, and height less than or equal to  $k+1$ .

*Proof.* Given  $T \in \mathcal{T}_{n,j}^{(k)}$ , label the edges by a postorder traversal. Read the labels by a preorder traversal to obtain a word  $\sigma \in S_n$ . Set  $\phi(T) = \sigma$ .

For example, consider the ordered tree  $T$  in Figure 3. When we read the labels by a

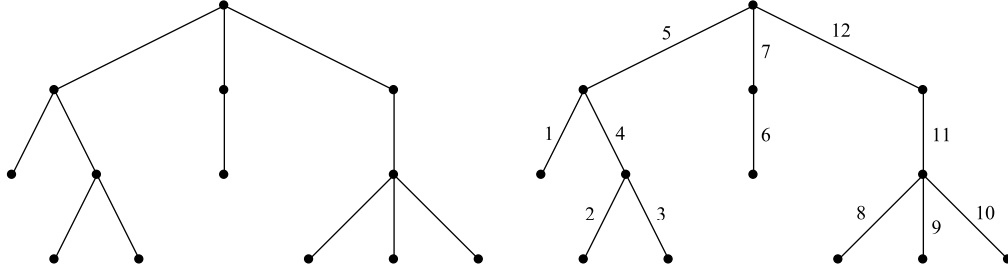


Figure 3: An ordered tree  $T$  (left),  $T$  with edges labeled by a postorder traversal (right).

preorder traversal, we obtain the permutation  $\sigma$  (which we write in two-line notation)

$$\sigma = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 \\ 5 & 1 & 4 & 2 & 3 & 7 & 6 & 12 & 11 & 8 & 9 & 10 \end{bmatrix}$$

Since Micheli and Rossin showed that  $\sigma \in S_n(231)$  and  $\phi$  is a bijection (see [12]), our Theorem is proved if we can show that  $\text{des}(\sigma) = j - 1$  and  $\text{maxdrop}(\sigma) \leq k - 1$ .

First we show that  $\text{des}(\sigma) = j - 1$ . Given any edge of  $T$ , let  $\sigma_i$  be its label from the postorder traversal, and let  $x$  be the vertex at the bottom of this edge.

If  $x$  is an internal node, then  $\sigma_{i+1}$  is the label on the leftmost edge immediately below  $x$ . Since the labeling is done with a postorder traversal, we have  $\sigma_i > \sigma_{i+1}$ .

If  $x$  is not an internal node (i.e. a leaf), then there is a vertex  $y$  with subtrees  $Y_1$  and  $Y_2$  such that  $\sigma_i$  is a label on an edge of  $Y_1$ ,  $\sigma_{i+1}$  is a label on an edge of  $Y_2$ , and  $Y_1$  is to the left of  $Y_2$ . It follows that  $\sigma_i < \sigma_{i+1}$ .

Since every vertex other than the root is at the bottom of a unique edge,  $\sigma$  has  $j - 1$  descents.

Next we show that  $\text{maxdrop}(\sigma) \leq k - 1$ . Suppose  $\sigma_i < i$ , and let  $x$  be the vertex at the bottom of the edge labeled  $\sigma_i$ .

If  $x$  is an internal node, then  $\sigma_i > \sigma_{i+1}$  as noted above, thus there is a larger drop size at position  $i + 1$  in  $\sigma$ . Since we want to find the maximum drop size, we need not consider the case that  $x$  is an internal node.

Now assume that  $x$  is not an internal node, and let  $m$  be the level of  $x$ . On the path from  $x$  to the root, there are  $m$  (possibly empty) subtrees along the left side of the path, as in Figure 4. Let  $|T_r|$  denote the number of edges in a tree  $T_r$ . Then we have

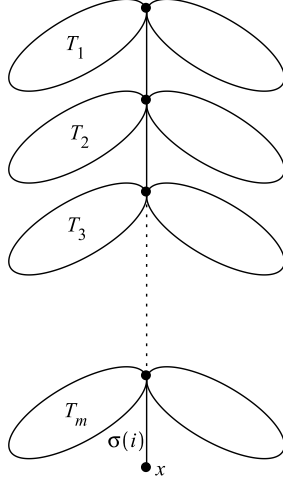


Figure 4: An ordered tree with edge labeled  $\sigma_i$  directly above a leaf  $x$  at level  $m$ .

$$|T_1| + |T_2| + \cdots + |T_m| = \sigma_i - 1,$$

since the edges in the subtrees  $T_1, T_2, \dots, T_m$  are precisely the edges which precede the edge labeled  $\sigma_i$  in the postorder traversal. The edges in the subtrees  $T_1, T_2, \dots, T_m$  along with the edges in the path from  $x$  to the root are precisely the edges which precede the edge labeled  $\sigma_i$  in the preorder traversal, therefore

$$i = |T_1| + |T_2| + \cdots + |T_m| + m = \sigma_i + m - 1.$$

Thus  $\sigma$  has a drop of size  $i - \sigma_i = m - 1$  at position  $i$ . Since  $m \leq k$ , we have  $\text{maxdrop}(\sigma) \leq k - 1$ .

□

The set of trees in  $\mathcal{T}_{n,j}^{(k)}$  are also in bijection with certain Dyck paths. A Dyck path of length  $2n$  is a path in the plane that starts at the point  $(0, 0)$  and ends at the point  $(2n, 0)$ . The path may consist only of up-steps  $(1, 1)$  and down-steps  $(1, -1)$ , and the path always stays on or above the  $x$ -axis. Let  $\mathcal{D}_{2n}$  denote the set of Dyck paths of length  $2n$ . Next we describe a couple statistics for Dyck paths. The height of a Dyck path is the highest  $y$ -coordinate attained in the path. A peak is a point in a Dyck path which is immediately preceded by an up-step, and immediately followed by a down-step. Let  $\mathcal{D}_{2n,j}^{(k)}$  denote the set of Dyck paths of length  $2n$  with  $j$  peaks and height less than or equal to  $k$ . The standard bijection from ordered trees to Dyck paths preserves height, and converts each leaf to a

peak. A tree with  $n$  edges and  $j + 1$  internal nodes has  $n + 1$  total nodes, thus  $n - j$  leaves. From this it follows that  $|\mathcal{T}_{n,j+1}^{(k)}| = |\mathcal{D}_{2n,n-j}^{(k)}|$ .

Next we provide a direct bijection from permutations in  $\mathcal{B}_{n,231,j}^{(k)}$  to Dyck paths in  $\mathcal{D}_{2n,n-j}^{(k+1)}$ . However, in subsequent sections of this paper we choose to use ordered trees to obtain enumeration results for the numbers  $a_{n,231,j}^{(k)}$  (and  $e_{n,231,j}^{(k)}$ ).

**Theorem 7.** *For all  $n \geq 1$  and all  $j, k \geq 0$ , there is a bijection  $\phi_n : \mathcal{B}_{n,231,j}^{(k)} \rightarrow \mathcal{D}_{2n,n-j}^{(k+1)}$ . In other words,  $a_{n,231,j}^{(k)}$  is equal to the number of Dyck paths of length  $2n$  with  $n - j$  peaks, and height less than or equal to  $k + 1$ .*

*Proof.* First we need to define the lifting of a path  $P \in \mathcal{D}_{2n}$  to path  $L(P) \in \mathcal{D}_{2n+2}$ . Let  $P = (p_1, \dots, p_{2n})$  where each  $p_i$  is either an up-step or a down-step. Then  $L(P)$  is obtained from  $P$  by appending an up-step at the start of  $P$ , and a down-step at the end of  $P$ . That is,  $L(P) = ((1, 1), p_1, \dots, p_{2n}, (1, -1))$ . An example is shown in Figure 5. Also, if  $P_1 \in \mathcal{D}_{2n}$  and  $P_2 \in \mathcal{D}_{2k}$ , then we let  $P_1 P_2 \in \mathcal{D}_{2n+2k}$  denote the path which starts with  $P_1$  followed by  $P_2$ .

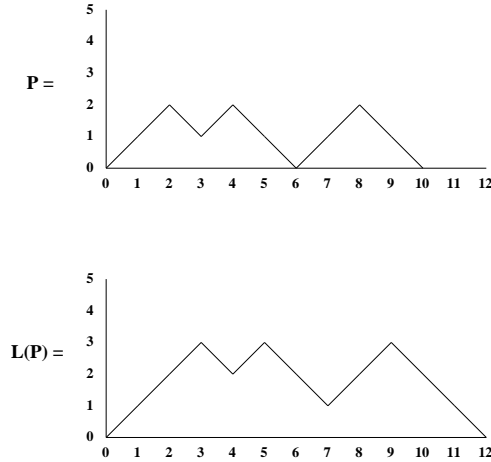


Figure 5: The lifting of a Dyck path.

For  $n = 1$ , we simply define  $\phi_1(\sigma) = ((1, 1), (1, -1))$ , i.e. and up-step followed by a down-step. For  $n > 1$  we define  $\phi_n$  recursively by cases as follows.

**Case 1.**  $\sigma_n = n$ .

Then  $\phi_n(\sigma) = P_1 P_2$  where  $P_1 = \phi_{n-1}(\sigma_1 \dots \sigma_{n-1})$  and  $P_2 = ((1, 1), (1, -1))$ .

**Case 2.**  $\sigma_1 = n$ .

Then  $\phi_n(\sigma) = L(\phi_{n-1}(\sigma_2 \dots \sigma_n))$ .

**Case 3.**  $\sigma_i = n$  where  $1 < i < n$ . In this case,  $\phi_n(\sigma) = P_1 P_2$  where  $P_1 = \phi_{i-1}(\text{red}(\sigma_1 \dots \sigma_{i-1}))$  and  $P_2 = L(\phi_{n-i}(\text{red}(\sigma_{i+1} \dots \sigma_n)))$ .

The first few values of this map are pictured in Figure 6, where  $\sigma$  is on the left and  $\phi_n(\sigma)$  is on the right.

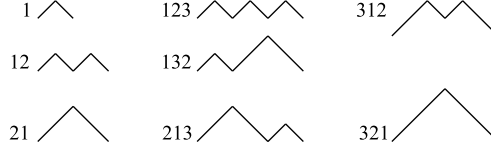


Figure 6: Values of  $\phi_n$  up to  $n = 3$ .

Since  $a_{n,231,j}^{(k)} = |\mathcal{T}_{n,j+1}^{(k+1)}| = |\mathcal{D}_{2n,n-j}^{(k+1)}|$ , it suffices to show that  $\phi_n$  is well-defined and injective. We induct on  $n$ . The base case  $n = 1$  is obvious. Now let  $n > 1$  and assume the theorem holds for all  $m < n$ . Let  $\sigma = \sigma_1 \dots \sigma_n \in \mathcal{B}_{n,231,j}^{(k)}$ , and let  $\sigma_i = n$ . Since  $\sigma$  avoids the pattern 231, we have  $\sigma_1 \dots \sigma_{i-1} \in \mathcal{B}_{i-1,231}^{(k)}$ , and  $\text{red}(\sigma_{i+1} \dots \sigma_n) \in \mathcal{B}_{n-i,231}^{(k-1)}$  (see also the proof of Theorem 1). To show  $\phi_n$  is well-defined, we consider the three cases for  $i$  in the definition of  $\phi_n$ .

**Case 1.**  $\sigma_n = n$ .

In this case,  $\phi_n(\sigma) = P_1 P_2$  where  $P_1 = \phi_{n-1}(\sigma_1 \dots \sigma_{n-1})$  and  $P_2 = ((1, 1), (1, -1))$ . Since  $n - 1 \notin \text{DES}(\sigma)$ , we have  $\sigma_1 \dots \sigma_{n-1} \in \mathcal{B}_{n-1,231,j}^{(k)}$ . By the inductive hypothesis we have  $P_1 \in \mathcal{D}_{2n-2,n-j-1}^{(k+1)}$ . Since appending  $P_2$  to  $P_1$  increases the length by two, increases the number of peaks by one, and does not change the height, it follows that  $P_1 P_2 \in \mathcal{D}_{n,n-j}^{(k+1)}$ .

**Case 2.**  $\sigma_1 = n$ .

In this case,  $\phi_n(\sigma) = L(\phi_{n-1}(\sigma_2 \dots \sigma_n))$ . Since  $1 \in \text{DES}(\sigma)$ , we have  $\sigma_2 \dots \sigma_n \in \mathcal{B}_{n-1,231,j-1}^{(k-1)}$ . By induction,  $\phi_{n-1}(\sigma_2 \dots \sigma_n) \in \mathcal{D}_{2n-2,n-j}^{(k)}$ . Since lifting a path increases the height by one, increasing the length by two, and adds no peaks, it follows that  $L(\phi_{n-1}(\sigma_2 \dots \sigma_n)) \in \mathcal{D}_{2n,n-j}^{(k+1)}$ .

**Case 3.**  $\sigma_i = n$  where  $1 < i < n$ .

In this case,  $\phi_n(\sigma) = P_1 P_2$  where  $P_1 = \phi_{i-1}(\text{red}(\sigma_1 \dots \sigma_{i-1}))$  and  $P_2 = L(\phi_{n-i}(\text{red}(\sigma_{i+1} \dots \sigma_n)))$ . Note that  $\text{red}(\sigma_1 \dots \sigma_{i-1}) = \sigma_1 \dots \sigma_{i-1}$ . Since  $i \in \text{DES}(\sigma)$  we have  $\sigma_1 \dots \sigma_{i-1} \in \mathcal{B}_{i-1,231,j_1}^{(k)}$ , and  $\text{red}(\sigma_{i+1} \dots \sigma_n) \in \mathcal{B}_{n-i,231,j_2}^{(k-1)}$ , where  $j_1 + j_2 + 1 = j$ . Then  $\phi_{i-1}(\text{red}(\sigma_1 \dots \sigma_{i-1})) \in \mathcal{D}_{2i-2,i-1-j_1}^{(k)}$ , and  $P_2 = L(\phi_{n-i}(\text{red}(\sigma_{i+1} \dots \sigma_n))) \in \mathcal{D}_{2n-2i+2,n-i-j_2}^{(k+1)}$ . Also,  $P_1 \in \mathcal{D}_{2i-2,i-1-j_1}^{(k+1)}$ . It follows that  $P_1 P_2 \in \mathcal{D}_{2n,n-j_1-j_2-1}^{(k+1)}$  as desired since  $n - j_1 - j_2 - 1 = n - j$ .

This proves  $\phi_n$  is well-defined.



To prove injectivity let  $\sigma, \pi \in \mathcal{B}_{n,231,j}^{(k)}$ , and suppose  $\sigma \neq \pi$ . If  $\sigma_i = \pi_i = n$  for some  $i$ , then for each case that  $i$  falls into in the definition of  $\phi_n$ , one can easily use the inductive hypothesis to prove that  $\phi_n(\sigma) \neq \phi_n(\pi)$ . So assume  $\sigma_{i_1} = n$  and  $\pi_{i_2} = n$  where  $i_1 \neq i_2$ . We consider the three possible combinations for  $i_1$  and  $i_2$  in the definition of  $\phi_n$ .

**Case I.**  $\sigma_n = n$  and  $\pi_1 = n$ .

Since  $\phi_n(\pi) = L(\phi_{n-1}(\pi_2 \dots \pi_n))$ , it follows that the second to last step of  $\phi_n(\pi)$  is a down-step. Then  $\phi_n(\sigma) \neq \phi_n(\pi)$  since the second to last step of  $\phi_n(\sigma)$  is an up-step.

**Case II.**  $\sigma_n = n$  and  $\pi_{i_2} = n$  where  $1 < i_2 < n$ .

Since  $\phi_n(\pi) = P_1 P_2$  where  $P_2 = L(\phi_{n-i_2}(\text{red}(\pi_{i_2+1} \dots \pi_n)))$ , we again have that the second to last step of  $\phi_n(\pi)$  is a down-step, whereas the second to last step of  $\phi_n(\sigma)$  is an up-step. Thus  $\phi_n(\sigma) \neq \phi_n(\pi)$ .

**Case III.**  $\sigma_1 = n$  and  $\pi_{i_2} = n$  where  $1 < i_2 < n$ .

In this case we note that  $\phi_n(\sigma) = L(\phi_{n-1}(\sigma_2 \dots \sigma_n))$ , so that the only points where  $\phi_n(\sigma)$  touches the  $x$ -axis are at  $(0, 0)$  and  $(2n, 0)$ . In contrast,  $\phi_n(\pi) = P_1 P_2$  where  $P_1 = \phi_{i_2-1}(\text{red}(\pi_1 \dots \pi_{i_2-1}))$ , so that  $\phi_n(\pi)$  must touch the  $x$ -axis at the point  $(2i_2 - 2, 0)$ . Thus  $\phi_n(\sigma) \neq \phi_n(\pi)$ . □

## 4 Recursions and closed form expressions for $a_{n,231,j}^{(k)}$

In this section we prove some recursions and closed form expressions for  $a_{n,231,j}^{(k)}$  and  $e_{n,231,j}^{(k)}$ . For certain cases of the values of  $n, j, k$ , we can find nice closed form expressions for these numbers. For the general case, it turns out that there is a closed form expression due to Kemp (see [5]) for a class of trees very closely related to  $\mathcal{T}_{n,j}^{(k)}$ . This formula can easily be translated to a closed form expression for  $a_{n,231,j}^{(k)}$ . We also find a recurrence for  $a_{n,231,j}^{(k)}$ . We conclude this section by showing that this recurrence leads to closed form expression for  $a_{n,231,j}^{(k)}$  which looks completely different from the formula due to Kemp.

**Theorem 8** ([5, Theorem 1]). *Let  $h_k(n, j)$  be the number of ordered trees with  $n$  nodes,  $j$  leaves, and height <sup>1</sup> less than or equal to  $k - 1$ . Then  $h_k(1, j) = \delta_{j,1}$ ,  $h_1(n, j) = \delta_{n,j} \delta_{n,1}$  where  $\delta_{n,j}$  is Kronecker's symbol. For  $k \geq 2$  and  $n \geq 2$*

$$h_k(n, j) = N(n - 1, j) - [Q_1(n, j, k) - 2Q_0(n, j, k) + Q_{-1}(n, j, k)],$$

where

$$Q_a(n, j, k) = \sum_{s \geq 1} \binom{n - s(k - 1) - 2}{j + s + a - 1} \binom{n + s(k - 1) - 2}{j - s - a - 1},$$

---

<sup>1</sup>In [5], the author uses the convention that the root is a vertex at level one, so we have translated this result to coincide with our definition of height.

and  $N(n, j)$  are the Narayana numbers given by

$$N(n, j) = \frac{1}{n} \binom{n}{j} \binom{n}{j-1}.$$

**Corollary 1.** *For all  $n \geq 1$ , and  $j, k \geq 0$  we have*

$$a_{n,231,j}^{(k)} = h_{k+2}(n+1, n-j).$$

*Proof.* A tree with  $n$  edges has  $n+1$  nodes. An ordered tree with  $n$  edges and  $j+1$  internal nodes has  $n+1 - (j+1) = n-j$  leaves. Thus from Theorem 6 we have

$$a_{n,231,j}^{(k)} = \left| \mathcal{T}_{n,j+1}^{(k+1)} \right| = h_{k+2}(n+1, n-j).$$

□

The Narayana numbers  $N(n, j)$  appear in several combinatorial problems (see A001263 in the OIES [13]). One interpretation is that  $N(n, j)$  is equal to the number of Dyck paths of length  $2n$  with  $j$  peaks. Another interpretation is that  $N(n, j)$  is equal to the number of ordered trees with  $n$  edges and  $j$  leaves. Next we show that  $a_{n,231,j}^{(k)}$  reduces to a Narayana number whenever  $k \geq j$ , extending the results from Theorem 3 and Theorem 4.

**Corollary 2.** *For all  $n, j \geq 1$ , and for all  $k \geq j$  we have*

$$a_{n,231,j}^{(k)} = N(n, n-j) = \frac{1}{n} \binom{n}{j} \binom{n}{j+1}.$$

*Proof.* An ordered tree with  $n$  edges and  $j+1$  internal nodes has height less than or equal to  $j+1$ , and  $n-j$  leaves. Thus whenever  $k \geq j$  we have

$$a_{n,231,j}^{(k)} = \left| \mathcal{T}_{n,j+1}^{(k+1)} \right| = N(n, n-j) = \frac{1}{n} \binom{n}{n-j} \binom{n}{n-j-1} = \frac{1}{n} \binom{n}{j} \binom{n}{j+1}.$$

Note that in general  $N(n, j) = N(n, n-j+1)$ , i.e. the Narayana numbers are symmetric, and this follows from the symmetry of the binomial coefficients.

□

In particular,  $a_{n,231,1}^{(k)} = \binom{n}{2}$  for  $k \geq 1$ , and  $a_{n,231,2}^{(k)} = \frac{1}{n} \binom{n}{2} \binom{n}{3} = \frac{(n-1)^2((n-1)^2-1)}{12}$  for  $k \geq 2$ , as expected from Theorem 3 and Theorem 4.

*Remark 1.* Corollary 2 also follows from Corollary 1 and Theorem 8 by noting that  $Q_a(n+1, n-j, k+2) = 0$  for  $a = -1, 0, 1$  whenever

$$n+1 - (k+1) - 2 < n-j+1 + a - 1.$$

And this inequality holds whenever  $k \geq j$ . Thus for  $k \geq j$  we have

$$a_{n,231,j}^{(k)} = h_{k+2}(n+1, n-j) = N(n, n-j).$$

Let  $\mathcal{E}_{n,231,j}^{(k)}$  be the set of permutations  $\sigma \in S_n(231)$  with  $\text{des}(\sigma) = j$  and  $\text{maxdrop}(\sigma) = k$ . Thus  $\mathcal{E}_{n,231,j}^{(k)} = \mathcal{B}_{n,231,j}^{(k)} - \mathcal{B}_{n,231,j}^{(k-1)}$ , and  $|\mathcal{E}_{n,231,j}^{(k)}| = e_{n,231,j}^{(k)} = a_{n,231,j}^{(k)} - a_{n,231,j}^{(k-1)}$ . We can interpret  $e_{n,231,j}^{(k)}$  as the number of ordered trees with  $n$  edges,  $j+1$  internal nodes, and height *equal* to  $k+1$ . For any ordered tree, the number of internal nodes is always less than or equal to its height. So  $e_{n,231,j}^{(k)} = 0$  if  $k \geq j+1$ . Using the tree interpretation, we will directly compute  $e_{n,231,j}^{(j)}$  by relating such trees to a certain set of weak compositions.

**Definition 1.** Let  $W_n(i, j)$  be the set of weak compositions  $(p_1, p_2, \dots, p_i; m_1, m_2, \dots, m_j)$  in  $\mathbb{N}^{i+j}$  such that (i)  $p_r \geq 1$  for  $r = 1, 2, \dots, i$ , (ii)  $m_r \geq 0$  for  $r = 1, 2, \dots, j$ , and (ii)  $\left(\sum_{r=1}^i p_r\right) + \left(\sum_{r=1}^j m_r\right) = n$ . In other words,  $W_n(i, j)$  is the set of weak compositions of  $n$  with  $i+j$  parts where the first  $i$  parts are positive.

**Proposition 1.** For all  $n, i, j \geq 0$  we have  $|W_n(i, j)| = \binom{n+j-1}{i+j-1}$ .

*Proof.* Let  $p'_r = p_r - 1$  for  $r = 1, 2, \dots, i$ . Then

$$(p_1, p_2, \dots, p_i; m_1, m_2, \dots, m_j) \in W_n(i, j)$$

if and only if

$$(p'_1, p'_2, \dots, p'_i; m_1, m_2, \dots, m_j) \in W_{n-i}(0, i+j).$$

$W_n(0, k)$  is simply the number of weak compositions of  $n$  into  $k$  parts, and  $|W_n(0, k)| = \binom{n+k-1}{k-1}$  (see [14]). Thus

$$|W_n(i, j)| = |W_{n-i}(0, i+j)| = \binom{n+j-1}{i+j-1}.$$

□

**Theorem 9.** For all  $j \geq 1$  we have

$$e_{n,231,j}^{(j)} = \binom{n+j-1}{2j}.$$

Consequently,

$$a_{n,231,j}^{(j-1)} = \frac{1}{n} \binom{n}{j} \binom{n}{j+1} - \binom{n+j-1}{2j}.$$

*Proof.* An ordered tree with  $j+1$  internal nodes and height equal to  $j+1$  must be a tree of the form shown in Figure 7 where each subtree  $T_r$  has only one internal node (its root), and  $|T_r| \geq 0$  for  $r = 1, 2, \dots, 2j$ . The subtree  $U$  must also have only one internal node, but  $|U| \geq 1$  so that the whole tree has height  $j+1$ . In other words,  $T_r \in \mathcal{T}_{m_r,1}^{(1)}$  with  $m_r \geq 0$  for  $r = 1, 2, \dots, 2j$ , and  $U \in \mathcal{T}_{p,1}^{(1)}$  with  $p \geq 1$ . Note that  $|\mathcal{T}_{m,1}^{(1)}| = 1$  for all  $m \geq 0$ , so every tuple  $(p; m_1, m_2, \dots, m_{2j}) \in W_{n-j}(1, 2j)$  corresponds to a unique tree with  $n$  edges,  $j+1$  internal nodes, and height equal to  $j+1$ . Thus by Proposition 1 we have

$$e_{n,231,j}^{(j)} = |W_{n-j}(1, 2j)| = \binom{n+j-1}{2j}.$$

□

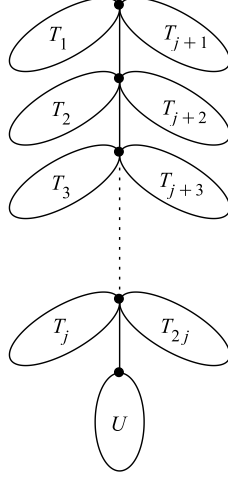


Figure 7: An ordered tree with  $j + 1$  internal nodes and height equal to  $j + 1$ .

In particular,  $e_{n,231,2}^{(2)} = \binom{n+1}{4}$  as expected from Theorem 5.

Since  $e_{n,231,j}^{(j)}$  is equal to a binomial coefficient, we also provide a direct bijection between such sets and permutations.

**Proposition 2.** *For all  $j \geq 1$  there is a bijection*

$$\phi : \binom{[n+j-1]}{2j} \rightarrow \mathcal{E}_{n,231,j}^{(j)},$$

where  $\binom{[n]}{j}$  is the set of  $j$ -element subsets of  $[n]$ .

*Proof.* Let  $S = \{a_1, a_2, \dots, a_{2j}\} \in \binom{[n+j-1]}{2j}$  with  $a_1 < a_2 < \dots < a_{2j}$ . We construct from  $S$  a sequence of nested intervals as follows. Let  $I_1^S = [a_1, b_1]$  where  $b_1 = \max\{S \cap [n]\}$ . Then for  $m = 2, 3, \dots, j$ , let  $I_m^S = [a_m, b_m]$  where

$$b_m = \begin{cases} b_{m-1} & \text{if } n + m - 1 \in S \\ \max\{[1, b_{m-1} - 1] \cap S\} & \text{otherwise} \end{cases}.$$

By construction we have  $a_1 < a_2 < \dots < a_j < b_j \leq b_{j-1} \leq b_{j-2} \leq \dots \leq b_1$ . Next define a map  $c_m^S$  which acts on permutations by cyclically rotating the letters in positions  $a_m, a_m + 1, \dots, b_m$ . More precisely, given  $\pi \in S_n$  let  $c_m^S(\pi) = \tau_1 \dots \tau_n$  where  $\tau_i = \pi_i$  if  $i \notin I_m^S$ , and  $\tau_{a_m} \tau_{a_m+1} \dots \tau_{b_m} = \pi_{b_m} \pi_{a_m} \pi_{a_m+1} \dots \pi_{b_m-1}$ . We then define

$$\phi(S) = c_j^S \circ c_{j-1}^S \circ \dots \circ c_1^S(\text{id}).$$

For example let  $n = 7$ ,  $j = 3$ , and let  $S = \{1, 3, 4, 5, 6, 8\} \in \binom{[9]}{6}$ . Then  $b_1 = \max\{S \cap [7]\} = 6$ , so  $I_1^S = [1, 6]$ . Next we find  $b_2$ . Since  $n + 2 - 1 = 8 \in S$ , we set  $b_2 = b_1 = 6$  and

$I_2^S = [3, 6]$ . Next we find  $b_3$ . Since  $n + 3 - 1 = 9 \notin S$ , we set  $b_3 = \max\{[1, b_2 - 1] \cap S\} = 5$  and  $I_3^S = [4, 5]$ . We can visualize the sequence of maps  $c_j^S \circ \dots \circ c_1^S$  by starting with the identity permutation and underlining the letters to be rotated in the next step.

$$\begin{aligned} \text{id} &= \underline{1} \ \underline{2} \ \underline{3} \ \underline{4} \ \underline{5} \ \underline{6} \ 7 \\ c_1^S(\text{id}) &= 6 \ 1 \ \underline{2} \ \underline{3} \ \underline{4} \ \underline{5} \ 7 \\ c_2^S \circ c_1^S(\text{id}) &= 6 \ 1 \ 5 \ \underline{2} \ \underline{3} \ 4 \ 7 \\ \phi(S) = c_3^S \circ c_2^S \circ c_1^S(\text{id}) &= 6 \ 1 \ 5 \ 3 \ 2 \ 4 \ 7 \end{aligned}$$

Since we already know that  $\left| \binom{[n+j-1]}{2j} \right| = \left| \mathcal{E}_{n,231,j}^{(j)} \right|$ , it suffices to show that  $\phi$  is well-defined and injective. To show  $\phi$  is well-defined, we induct on  $j$ . We want to show that

- (I)  $\phi(S) = \mathcal{E}_{n,231,j}^{(j)}$  whenever  $S = \{a_1 < a_2 < \dots < a_{2j}\} \in \binom{[n+j-1]}{2j}$ ,
- (II)  $\pi_{a_j} > \pi_{b_j}$ , where  $\phi(S) = \pi_1 \dots \pi_n$ , and  $b_1, \dots, b_j$  are as described in the definition of  $\phi$ ,
- (III)  $i - \pi_i = j$  for  $i = a_j + 1, a_j + 2, \dots, b_j$ .

The base case is obvious (and coincides with the bijection given in Theorem 2 when  $j = 1$ ). Now let  $j > 1$  and assume the result holds for all  $k < j$ . Let  $S = \{a_1 < a_2 < \dots < a_{2j}\}$ , and let

$$b'_j = \begin{cases} b_j & \text{if } b_j \in [n] \\ n + j - 1 & \text{otherwise} \end{cases}.$$

It follows that  $\phi(S) = c_j^S(\phi(T))$ , where  $T = S - \{a_j, b'_j\}$ . Let  $\sigma = \phi(T)$ , and let  $\pi = \phi(S)$ . Thus  $\pi_i = \sigma_i$  if  $i \notin [a_j, b_j]$ , and  $\pi_{a_j} \pi_{a_j+1} \dots \pi_{b_j} = \sigma_{b_j} \sigma_{a_j} \sigma_{a_j+1} \dots \sigma_{b_j-1}$ . By induction we have that  $i - \sigma_i = j - 1$  for  $i = a_{j-1} + 1, a_{j-1} + 2, \dots, b_{j-1}$ . Since  $a_{j-1} < a_j < b_j \leq b_{j-1}$ , then for  $i = a_j + 1, a_j + 2, \dots, b_j$  we have

$$i - \pi_i = i - \sigma_{i-1} = i - (i - j) = j,$$

which proves (III). Next we check the drop size of  $\pi$  at position  $a_j$

$$a_j - \pi_{a_j} = a_j - \sigma_{b_j} = a_j - (b_j - j + 1) < b_j - (b_j - j + 1) = j - 1,$$

thus  $\text{maxdrop}(\pi) = j$ . (II) follows immediately since

$$\pi_{a_j} = \sigma_{b_j} = b_j - j + 1 > b_j - j = \pi_{b_j}.$$

Next we prove the claim that  $\text{des}(\pi) = j$ . By induction we know that  $\text{des}(\sigma) = j - 1$  and  $\sigma_{a_j} < \sigma_{a_j+1} < \dots < \sigma_{b_j}$ . Moreover, (II) and (III) applied to  $\pi$  shows us that  $\text{DES}(\pi) \cap [a_j, b_j - 1] = \{a_j\}$ . Since  $\pi_{b_j} = b_j - j = \sigma_{b_j} - 1$  and  $\pi_{b_j+1} = \sigma_{b_j+1}$ , it follows that  $\pi_{b_j}, \pi_{b_j+1}$

have the same relative order as  $\sigma_{b_j}, \sigma_{b_j+1}$ . Thus  $b_j \in \text{DES}(\pi)$  if and only if  $b_j \in \text{DES}(\sigma)$ . If  $a_j - 1 \notin \text{DES}(\sigma)$ , then

$$\pi_{a_j-1} = \sigma_{a_j-1} < \sigma_{a_j} < \sigma_{b_j} = \pi_{a_j},$$

which implies  $a_j - 1 \notin \text{DES}(\pi)$ . Conversely, suppose that  $a_j - 1 \in \text{DES}(\sigma)$ . Since  $\text{DES}(\sigma) \cap [a_{j-1} + 1, b_{j-1}] = \emptyset$ , we have  $a_j - 1 \leq a_{j-1}$ . But since  $a_{j-1} \leq a_j - 1$ , we have  $a_j - 1 = a_{j-1}$ . Then

$$\pi_{a_j-1} = \sigma_{a_j-1} = \sigma_{a_{j-1}} > \sigma_{b_{j-1}} \geq \sigma_{b_j} = \pi_{a_j},$$

which implies  $a_j - 1 \in \text{DES}(\pi)$ . Since  $\pi_i = \sigma_i$  for all  $i \notin [a_j, b_j]$ , it follows that  $\text{DES}(\pi) = \text{DES}(\sigma) \uplus \{a_j\}$ , thus  $\text{des}(\pi) = \text{des}(\sigma) + 1 = j$  which proves (I).

Next we prove that  $\phi$  is injective. Let  $S = \{a_1 < a_2 < \dots < a_{2j}\} \in \binom{[n+j-1]}{2j}$  from which we construct the sequence of intervals  $I_1^S \supset I_2^S \supset \dots \supset I_j^S$  with  $I_m^S = [a_m, b_m]$  for  $m = 1, \dots, j$ . And let  $U = \{x_1 < x_2 < \dots < x_{2j}\} \in \binom{[n+j-1]}{2j}$  from which we construct the sequence of intervals  $I_1^U \supset I_2^U \supset \dots \supset I_j^U$  with  $I_m^U = [x_m, y_m]$  for  $m = 1, \dots, j$ . Suppose  $S \neq U$ , then we claim there is at least one index  $m$  such that  $I_m^S \neq I_m^U$ . Suppose there is some  $a_k \in [n] \cap S$  such that  $a_k \notin U$ . Then there is at least one interval  $I_m^S$  such that  $a_m = a_k$  or  $b_m = a_k$ , but there is no interval  $I_m^U$  with  $x_m = a_k$  or  $y_m = a_k$ . Next suppose that for some  $m$  with  $2 \leq m \leq j$  we have  $n + m - 1 \in S$ , but  $n + m - 1 \notin U$ . Then  $b_m = b_{m-1}$  and  $y_m \neq y_{m-1}$ . It follows that either  $I_m^S \neq I_m^U$  or  $I_{m-1}^S \neq I_{m-1}^U$ . This proves the claim.

Let  $\pi = \phi(S)$  and  $\sigma = \phi(U)$ . It follows from (II) and (III) above and the fact that  $a_1 < a_2 < \dots < a_j$ , that  $\text{DES}(\pi) = \{a_1, a_2, \dots, a_j\}$ . Similarly,  $\text{DES}(\sigma) = \{x_1, x_2, \dots, x_j\}$ . If  $a_m \neq x_m$  for some  $m = 1, 2, \dots, j$ , then  $\text{DES}(\pi) \neq \text{DES}(\sigma)$  and  $\pi \neq \sigma$ . It follows from (III) above and from the definition of  $\phi$  that  $i - \pi_i = j - k$  for  $i = b_{j-k+1} + 1, b_{j-k+1} + 2, \dots, b_{j-k}$  and  $k = 1, 2, \dots, j - 1$ . It is also clear that  $\pi_i = i$  for  $i > b_1$ . Now suppose that  $a_m \neq x_m$  for some  $m = j + 1, j + 2, \dots, 2j$ . This implies that  $b_m \neq y_m$ , so assume  $b_m > y_m$ . First we note that  $b_m - \pi_{b_m} \geq m$ , since  $\phi$  reduces the letter in position  $b_m$  by one at least  $m$  times (it will be reduced by one more than  $m$  times if  $b_{m-1} = b_m$ ). On the other hand since  $y_m < b_m$  we have  $b_m - \sigma_{b_m} < m$ , thus  $\pi \neq \sigma$ . This proves  $\phi$  is injective.  $\square$

A method similar to the one used to prove Theorem 9, proves the following theorem.

**Theorem 10.** *For all  $j \geq 2$  we have*

$$e_{n,231,j}^{(j-1)} = (2j - 3) \binom{n+j-2}{2j}.$$

Consequently,

$$a_{n,231,j}^{(j-2)} = \frac{1}{n} \binom{n}{j} \binom{n}{j+1} - \binom{n+j-1}{2j} - (2j - 3) \binom{n+j-2}{2j}.$$

*Proof.* We interpret  $e_{n,231,j}^{(j-1)}$  as the number of ordered trees with  $n$  edges,  $j + 1$  internal nodes, and height equal to  $j$ . Such trees have the form shown in Figure 8 where the

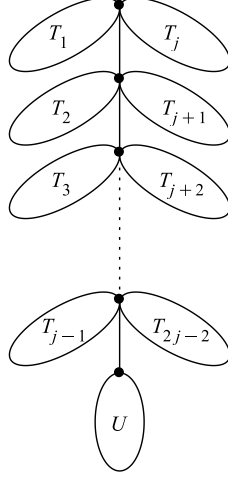


Figure 8: An ordered tree with  $j + 1$  internal nodes and height equal to  $j$ .

subtree  $U \in \mathcal{T}_{p,1}^{(1)}$  with  $p \geq 1$ , and the subtree  $T_{j-1} \in \mathcal{T}_{m_{j-1},1}^{(1)}$  with  $m_{j-1} \geq 0$ . The remaining subtrees  $T_r \in \mathcal{T}_{m_r,d_r}^{(2)}$  with  $m_r \geq 0$  for  $1 \leq r \leq 2j - 2$  with  $r \neq j - 1$ . The idea here is that the leftmost leaf in  $U$  is the first vertex (in postorder) to reach the full height  $j$ . We also require that exactly one of the subtrees  $T_r$  with  $r \neq j - 1$  has two internal nodes so that the resulting tree has  $j + 1$  total internal nodes, i.e.

$$\sum_{\substack{1 \leq r \leq 2j-2 \\ r \neq j-1}} d_r = 2j - 2.$$

To get a total of  $n$  edges, we also require that

$$p + \sum_{r=1}^{2j-2} m_r = n - j + 1.$$

Then every such choice of subtrees  $U, T_1, T_2, \dots, T_{2j-2}$  corresponds to a unique tree with  $n$  edges,  $j + 1$  internal nodes, and height equal to  $j$ . Since  $|\mathcal{T}_{m_r,d_r}^{(2)}| = a_{m_r,231,d_r-1}^{(1)} = \binom{m_r}{2d_r-2}$ .

Thus

$$\begin{aligned}
e_{n,231,j}^{(j-1)} &= \sum_{\substack{(p;m_1,\dots,m_{2j-2}) \in W_{n-j+1}(1,2j-2) \\ (d_1,d_2,\dots,d_{j-2},d_j,d_{j+1},\dots,d_{2j-2}) \in W_{2j-2}(2j-3,0)}} \prod_{r \neq j-1} \binom{m_r}{2d_r-2} \\
&= (2j-3) \left( \sum_{(p;m_1,\dots,m_{2j-2}) \in W_{n-j+1}(1,2j-2)} \binom{m_1}{2} \right) \\
&= (2j-3) \left( \sum_{(p;m_1,\dots,m_{2j-2}) \in W_{n-j}(0,2j-1)} \binom{m_1}{2} \right) \\
&= (2j-3) \sum_{m_1=0}^{n-j} \binom{m_1}{2} \sum_{(p;m_2,\dots,m_{2j-2}) \in W_{n-j-m_1}(0,2j-2)} 1 \\
&= (2j-3) \sum_{m_1=0}^{n-j} \binom{m_1}{2} |W_{n-j-m_1}(0,2j-2)| \\
&= (2j-3) \sum_{m_1=0}^{n-j} \binom{m_1}{2} \binom{n+j-m_1-3}{2j-3} \\
&= (2j-3) \sum_{m_1=0}^{n-j} \binom{m_1}{2} \binom{n+j-3-m_1}{(2j-1)-2} \\
&= (2j-3) \binom{n+j-2}{2j}.
\end{aligned}$$

The last step follows from the identity

$$\sum_{m=0}^n \binom{m}{j} \binom{n-m}{k-j} = \binom{n+1}{k+1},$$

which holds for all  $n \geq k \geq j \geq 0$ , and can easily be proved by induction.  $\square$

A similar method could be used to compute  $e_{n,231,j}^{(j-2)}$ . However, the proof continues to become more complicated. We hope the reader is convinced that the proof and resulting formula for  $e_{n,231,j}^{(j-2)}$  will be somewhat unpleasant, and that this method will become even more unpleasant as we continue to lower the maximum drop size. Instead, we will show later that  $a_{n,j}^{(k)}$  can be expressed as a (positive) sum of products of binomial coefficients.

Our next goal is to obtain a recurrence for  $a_{n,231,j}^{(k)}$ . We accomplish this using a bijection to find a recurrence for trees with  $n$  edges,  $j$  leaves, and height less than or equal to  $k$ . Let  $\mathcal{N}(n, j, k)$  denote the set of ordered trees with  $n$  edges,  $j$  leaves, and height less than or equal to  $k$ , and let  $N(n, j, k) = |\mathcal{N}(n, j, k)|$ . In other words,  $N(n, j, k)$  are the Narayana numbers refined by height. For convenience, we let  $\mathcal{N}(0, 0, k)$  be the set containing the tree



with one vertex and no edges, hence  $N(0, 0, k) = 1$  for all  $k \geq 0$ . Note that in terms of Dyck paths we have  $N(n, j, k) = |\mathcal{D}_{2n, j}^{(k)}|$ , i.e. the number of Dyck paths of length  $2n$  with  $j$  peaks and height less than or equal to  $k$ . We find a recurrence for  $N(n, j, k)$ , which can easily be translated into a recurrence for  $a_{n, 231, j}^{(k)}$ .

**Theorem 11.** *For all  $k \geq 1$  and for all  $n \geq j \geq 1$  we have*

$$N(n, j, k) = \sum_{i=0}^{n-j} N(n-j, i, k-1) \binom{2n-j-i}{2n-2j}. \quad (11)$$

*By replacing  $i$  with  $n-j-i$  we obtain*

$$N(n, j, k) = \sum_{i=0}^{n-j} N(n-j, n-j-i, k-1) \binom{n+i}{2n-2j}. \quad (12)$$

*Proof.* We construct a map

$$s : \bigcup_{i=0}^{n-j} \mathcal{N}(n-j, i, k-1) \times W_j(i, 2n-2j-i+1) \rightarrow \mathcal{N}(n, j, k),$$

which we will show is a bijection. Let  $T \in \mathcal{N}(n-j, i, k-1)$ , and let

$$c = (l_1, \dots, l_i; n_1, \dots, n_{2n-2j-i+1}) \in W_j(i, 2n-2j-i+1)$$

for some  $i$  such that  $0 \leq i \leq n-j$ . We describe  $s(T, c)$  via a composition of maps,  $s = s_1 \circ s_2 \circ \dots \circ s_k$ . Let  $U_h = s_h \circ s_{h+1} \circ \dots \circ s_k(T, c)$  with  $2 \leq h \leq k$ , and set  $U_{k+1} = (T, c)$ . Construct an ordered tree  $U_{h-1} = s_{h-1}(U_h)$  by visiting the vertices at level  $h-1$  of  $U_h$  from right to left, and (possibly) adding edges to each vertex  $x$  as follows:

- If  $x$  is the  $p^{\text{th}}$  leaf visited in the process of applying  $s_{h-1}, s_h, \dots, s_k$  to  $(T, c)$ , then attach the unique tree from  $\mathcal{N}(l_p, l_p, 1)$  as a subtree below  $x$ .
- If  $x$  is an internal node with degree say  $d$  (i.e.  $x$  has  $d$  children), then attach the tree from  $\mathcal{N}(n_r, n_r, 1)$  to the right of the rightmost edge below  $x$ . Here  $r = 1 + \sum (\deg(y) + 1)$  where the sum is over all internal nodes previously visited in the process of applying  $s_{h-1}, s_h, \dots, s_k$  to  $(T, c)$ . Then for  $m = 1, 2, \dots, d$ , attach the tree from  $\mathcal{N}(n_{r+m}, n_{r+m}, 1)$  to the left of the  $m$ -th rightmost edge below  $x$ .

Consider the following example. Let  $n = 14, j = 10, k = 3, i = 3$ . Let  $T \in \mathcal{N}(4, 3, 2)$  be the tree shown in Figure 9.

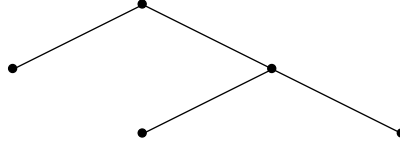


Figure 9: An ordered tree  $T \in \mathcal{N}(4, 3, 2)$ .

Let

$$c = (1, 2, 2; 1, 2, 0, 0, 1, 1) \in W_{10}(3, 6).$$

Find  $s_3(T, c)$  by visiting the vertices at level 2 from right to left. Both vertices are leaves, so we attach  $l_1 = 1$  edge to the right vertex, and  $l_2 = 2$  edges to the left vertex, where the dashed edges represent the added edges (see Figure 10).

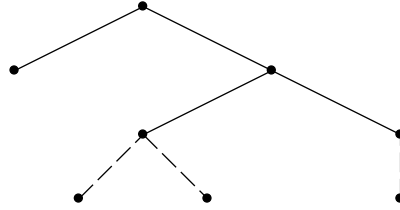


Figure 10:  $s_3(T, c)$ ,  $c = (\mathbf{1}, \mathbf{2}, 2; 1, 2, 0, 0, 1, 1)$ .

We continue by applying  $s_2$ . The rightmost vertex on level 1 is an internal node, so we attach  $n_1 = 1$  edge to the right of the rightmost edge, attach  $n_2 = 2$  edges in the middle, and attach  $n_3 = 0$  edges to the left of the leftmost edge. The next vertex on level 1 is a leaf, so we attach  $l_3 = 2$  edges below this vertex (see Figure 11).

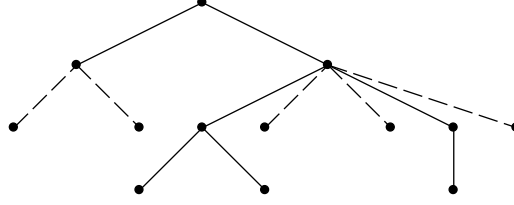


Figure 11:  $s_2 \circ s_3(T, c)$ ,  $c = (1, 2, \mathbf{2}; 1, \mathbf{2}, \mathbf{0}, 0, 1, 1)$ .

Lastly we apply  $s_1$  to obtain  $s(T, c) = s_1 \circ s_2 \circ s_3(T, c)$  (see Figure 12).

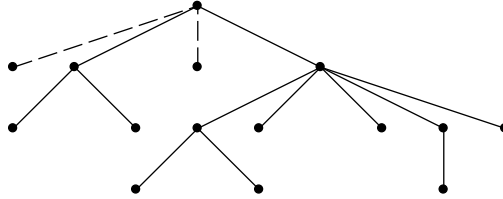


Figure 12:  $s_1 \circ s_2 \circ s_3(T, c)$ ,  $c = (1, 2, 2; 1, 2, 0, \mathbf{0}, \mathbf{1}, \mathbf{1})$ .

Next we show that  $s$  is well-defined. Note that  $T$  has  $n - j$  edges and  $c$  is a weak composition of  $j$ . Thus applying  $s$  will add  $j$  edges to  $T$ , so  $s(T, c)$  has  $n$  edges. Since  $l_1, \dots, l_i$  are all positive, every leaf of  $T$  has edges added to it, and is therefore not a leaf in  $s(T, c)$ . On the other hand, every edge added to  $T$  creates a leaf, so  $s(T, c)$  has  $j$  leaves. Since  $T$  has height less than or equal to  $k - 1$ , it is clear that  $s(T, c)$  has height less than or equal to  $k$ . Furthermore,  $T$  has  $i$  leaves and the first  $i$  parts of  $c$  are positive. We also need to check that  $c$  has the appropriate number of parts for adding edges to internal nodes.

This follows from the fact that

$$\begin{aligned}
\sum_{\substack{x \text{ is an internal} \\ \text{node of } T}} (1 + \deg(x)) &= |\{\text{internal nodes of } T\}| + \sum_{\substack{x \text{ is an internal} \\ \text{node of } T}} \deg(x) \\
&= |\{\text{vertices of } T\}| - |\{\text{leaves of } T\}| + |\{\text{edges of } T\}| \\
&= (n - j + 1) - i + (n - j) \\
&= 2n - 2j - i + 1
\end{aligned}$$

Next we describe the inverse map of  $s$ , which we denote by  $f$ . We have chosen the letter  $s$  to correspond to spring, since the tree "grows" edges during this map. And the letter  $f$  corresponds to fall since we will remove edges during this map. Let  $T \in \mathcal{N}(n, j, k)$ . Again we describe  $f(T)$  via a composition of maps  $f_k \circ f_{k-1} \circ \cdots \circ f_1$ . Let  $(V_h, c_h) = f_h \circ f_{h-1} \circ \cdots \circ f_1(T)$  with  $1 \leq h \leq k - 2$ , where  $V_h$  is an ordered tree and  $c_h$  is a weak composition, and let  $(V_0, c_0) = (T, \emptyset)$ . Construct  $(V_{h+1}, c_{h+1}) = f_{h+1}(V_h, c_h)$  by visiting the vertices at level  $h$  of  $V_h$  from *left to right*, and removing all single edges below each vertex. The weak composition  $c_{h+1}$  is obtained from  $c_h$  by recording at each vertex  $x$ , the numbers of edges removed as follows:

- If  $x$  has  $p$  children and all subtrees below  $x$  have height one (in other words  $x$  has no grandchildren), then append a  $p$  to the beginning of the positive parts of  $c_h$ .
- Suppose  $x$  has  $Y_1, Y_2, \dots, Y_d$  (from left to right) subtrees of height greater than or equal to 2. Let  $r_1$  be the number of single edges below  $x$  and to the left of  $Y_1$ . For  $m = 1, 2, \dots, d - 1$ , let  $r_m$  equal the number of single edges below  $x$  between  $Y_m$  and  $Y_{m+1}$ . Let  $r_{d+1}$  be the number of single edges below  $x$  and to the right of  $Y_d$ . Then append the parts  $(r_{d+1}, r_d, \dots, r_1)$  to the beginning of the nonnegative parts of  $c_h$ .

Here is an example, let  $T = V_0 \in \mathcal{N}(10, 7, 3)$  be the ordered tree in Figure 13.

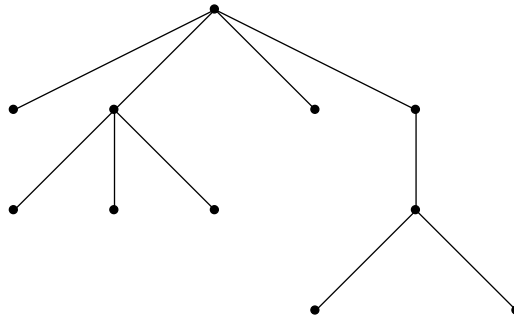


Figure 13: An ordered tree  $T = V_0$ , and set  $c_0 = \emptyset$ .

The root has 2 subtrees  $Y_1$  and  $Y_2$  with height greater than or equal to 2. There is  $n_1 = 1$  single edge to the left of  $Y_1$ , there is  $n_2 = 1$  single edge to the right of  $Y_1$ , and  $n_3 = 0$  single edges to the right of  $Y_2$ . We remove these single edges and record the number of edges removed as nonnegative parts of the weak composition  $c_1$ , i.e.  $c_1 = (n_3, n_2, n_1) = (0, 1, 1)$  (see Figure 14).

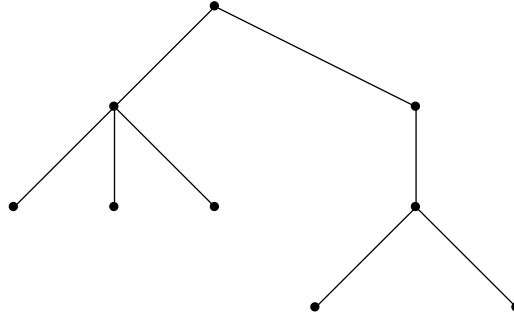


Figure 14: The ordered tree  $V_1$ , and  $c_1 = (0, 1, 1)$ .

Next we visit the vertices at level 1. The first vertex (moving from left to right) has only single edges. We remove these  $l_1 = 3$  edges and record the number of edges removed as a positive part of the weak composition  $c_2$ . The next vertex has one subtree of height 2. There are no single edges, so we record two zeros as nonnegative parts of the weak composition  $c_2$ , i.e.  $n_4 = n_5 = 0$  (see Figure 15).

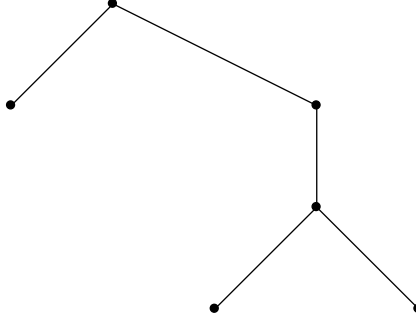


Figure 15: The ordered tree  $V_2$ , and  $c_2 = (3; 0, 0, 0, 1, 1)$ .

At level 2, there is one vertex with only single edges, so we remove them and record the number of edges removed as a positive part of the weak composition  $c_3$ , i.e.  $l_2 = 2$  (see Figure 16).

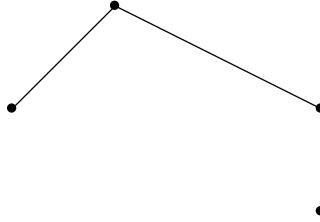


Figure 16: The ordered tree  $V_3$ , and  $c_3 = (2, 3; 0, 0, 0, 1, 1)$ .

Next we show  $f$  is well-defined. When applying  $f_h$ , we never visit a vertex which is a leaf since such a vertex would have been removed when applying  $f_{h-1}$ . Since we remove from  $T$  precisely all edges which have a leaf at the bottom, we see that  $V_k$  has  $n - j$  edges. Since  $V_k$  has  $n - j$  edges, the number of leaves of  $V_k$  is less than or equal to  $n - j$ . Clearly, the height of  $V_k$  is one less than the height of  $T$ . Thus  $V_k \in \mathcal{N}(n - j, i, k - 1)$  for some  $0 \leq i \leq n - j$ .

Since the map  $f$  removes  $j$  edges from  $T$ , and since  $c_k$  records the total number of edges removed,  $c_k$  is a weak composition of  $j$ . A leaf of  $V_k$  is created only when we visit a vertex with only single edges below. The number of such edges is recorded as a positive part in the weak composition  $c_k$ . So if  $V_k$  has  $i$  leaves, then  $c_k$  has  $i$  positive parts. Lastly, the total number of parts (positive and nonnegative) of  $c_k$  is given by

$$\begin{aligned} \sum_{\substack{x \text{ is a} \\ \text{vertex of } V_k}} (1 + \deg(x)) &= |\{\text{vertices of } V_k\}| + |\{\text{edges of } V_k\}| \\ &= (n - j + 1) + (n - j) \\ &= 2n - 2j + 1, \end{aligned}$$

thus  $c_k$  has  $2n - 2j + 1 - i$  nonnegative parts.

It is clear by construction that  $f$  is the inverse of  $s$ .

The Theorem now follows from the fact that (see Proposition 1)

$$|W_j(i, 2n - 2j - i + 1)| = \binom{2n - j - i}{2n - 2j}.$$

□

The following is a recurrence for the Narayana number due to Zabrocki (see A001263 in the OEIS [13]).

**Corollary 3.** *For  $n \geq j \geq 2$  we have*

$$N(n, j) = \sum_{i=1}^{j-1} N(j-1, i) \binom{n+i-1}{2j-2}.$$

*Proof.* Assume  $j \geq 2$ . If  $k \geq n$ , then  $N(n, j, k) = N(n, j)$  and  $N(n-j, n-j-i, k-1) = N(n-j, n-j-i)$  so that (12) becomes

$$N(n, j) = \sum_{i=0}^{n-j} N(n-j, n-j-i) \binom{n+i}{2n-2j}.$$

Since the Narayana number are symmetric,  $N(n, j) = N(n, n-j+1)$ , we have

$$\begin{aligned} N(n, j) &= N(n, n-j+1) \\ &= \sum_{i=0}^{j-1} N(j-1, j-1-i) \binom{n+i}{2j-2} \\ &= \sum_{i=0}^{j-2} N(j-1, i+1) \binom{n+i}{2j-2} \end{aligned}$$

The result now follows from replacing  $i$  by  $i-1$ .

□

Theorem 11 also gives us a recurrence for  $a_{n,231,j}^{(k)}$ .

**Corollary 4.** *For  $k \geq 1$  we have*

$$a_{n,231,j}^{(k)} = \sum_{i=0}^j a_{j,231,i}^{(k-1)} \binom{n+i}{2j},$$

where  $a_{0,231,0}^{(k-1)} := 1$  for all  $k \geq 1$ .

*Proof.*

$$\begin{aligned} a_{n,231,j}^{(k)} &= \left| \mathcal{T}_{n,j+1}^{(k+1)} \right| = N(n, n-j, k+1) = \sum_{i=0}^j N(j, j-i, k) \binom{n+i}{2j} \\ &= \sum_{i=0}^j a_{j,231,i}^{(k-1)} \binom{n+i}{2j}. \end{aligned}$$

□

The recurrence in Corollary 4 can be iterated to obtain a closed form expressions for  $a_{n,231,j}^{(k)}$ . Indeed,  $a_{j,231,i}^{(0)}$  is the number of permutations in  $S_j(231)$  with  $i$  descents and no drops. Since the identity permutation is the only permutation with no drops, we see that  $a_{j,231,i}^{(0)} = 1$  if  $i = 0$ , and  $a_{j,231,i}^{(0)}$  is zero otherwise. Hence

$$a_{n,231,j}^{(1)} = \sum_{i=0}^j a_{j,231,i}^{(0)} \binom{n+i}{2j} = \binom{n}{2j},$$

as expected (see Theorem 2). We iterate to obtain the following formula.

**Corollary 5.** *For all  $n, j \geq 0$  we have*

$$a_{n,231,j}^{(2)} = \sum_{j \geq i \geq 0} a_{j,231,i}^{(1)} \binom{n+i}{2j} = \sum_{j \geq i \geq 0} \binom{j}{2i} \binom{n+i}{2j}.$$

We continue iterating to obtains more formulas. In each case the formula holds for all  $n, j \geq 0$ .

$$\begin{aligned} a_{n,231,j}^{(3)} &= \sum_{j \geq i_2 \geq i_1 \geq 0} \binom{i_2}{2i_1} \binom{j+i_1}{2i_2} \binom{n+i_2}{2j}. \\ a_{n,231,j}^{(4)} &= \sum_{j \geq i_3 \geq i_2 \geq i_1 \geq 0} \binom{i_2}{2i_1} \binom{i_3+i_1}{2i_2} \binom{j+i_2}{2i_3} \binom{n+i_3}{2j}. \\ a_{n,231,j}^{(5)} &= \sum_{j \geq i_4 \geq i_3 \geq i_2 \geq i_1 \geq 0} \binom{i_2}{2i_1} \binom{i_3+i_1}{2i_2} \binom{i_4+i_2}{2i_3} \binom{j+i_3}{2i_4} \binom{n+i_4}{2j}. \end{aligned}$$

A pattern emerges, giving us a formula for  $a_{n,231,j}^{(k)}$ .



**Theorem 12.** For all  $n, j \geq 0$  and all  $k \geq 2$  we have

$$a_{n,231,j}^{(k)} = \sum_{j \geq i_{k-1} \geq \dots \geq i_1 \geq 0} \left( \prod_{m=0}^{k-1} \binom{i_{m+2} + i_m}{2i_{m+1}} \right),$$

where  $i_0 := 0$ ,  $i_k := j$ ,  $i_{k+1} := n$ .

*Proof.* Induct on  $k$ . The base case  $k = 2$  follows from Corollary 5. Now let  $k \geq 3$  and assume the result holds for  $k - 1$ . Then from Corollary 4 we have

$$a_{n,231,j}^{(k)} = \sum_{j \geq p_{k-1} \geq 0} a_{j,231,p_{k-1}}^{(k-1)} \binom{n + p_{k-1}}{2j}.$$

Use the induction hypothesis to substitute an expression for  $a_{j,231,p_{k-1}}^{(k-1)}$ .

$$a_{n,231,j}^{(k)} = \sum_{\substack{j \geq p_{k-1} \geq 0 \\ p_{k-1} \geq i_{k-2} \geq \dots \geq i_1 \geq 0}} \left( \prod_{m=0}^{k-4} \binom{i_{m+2} + i_m}{2i_{m+1}} \right) \binom{p_{k-1} + i_{k-3}}{2i_{k-2}} \binom{j + i_{k-2}}{2p_{k-1}} \binom{n + p_{k-1}}{2j}.$$

The result now follows from replacing  $p_{k-1}$  with  $i_{k-1}$ .

□

We translate Theorem 12 into an explicit formula for  $N(n, j, k)$  (which may be interpreted in terms of ordered trees, or in terms of Dyck paths).

**Corollary 6.** For  $n \geq j \geq 0$  and  $k \geq 3$  we have

$$N(n, j, k) = a_{n,231,n-j}^{(k-1)} = \sum_{n-j \geq i_{k-2} \geq \dots \geq i_1 \geq 0} \left( \prod_{m=0}^{k-2} \binom{i_{m+2} + i_m}{2i_{m+1}} \right),$$

where  $i_0 := 0$ ,  $i_{k-1} := n - j$ ,  $i_k := n$ .

## 5 Resulting Identities

In the previous section we proved that the set of permutations in  $S_n(231)$  with  $j$  descents and maximum drop less than or equal to  $k$ , is in bijective correspondence with the set of ordered trees with  $n$  edges,  $j + 1$  internal nodes, and height less than or equal to  $k + 1$ . We also found two seemingly different closed form expressions for the number of such trees: one due to Kemp [5] (Theorem 8 and Corollary 1), and another resulting from iterating our recurrence (Theorem 12). This leads to some remarkable identities.

**Theorem 13.** For  $n \geq 1$  and  $j \geq 0$  we have

$$a_{n,231,j}^{(1)} = h_3(n + 1, n - j). \tag{13}$$

Consequently

$$\binom{n}{2j} = N(n, j+1) - \left[ \tilde{Q}_1(n, j, 3) - 2\tilde{Q}_0(n, j, 3) + \tilde{Q}_{-1}(n, j, 3) \right], \quad (14)$$

where

$$N(n, j+1) = \frac{1}{n} \binom{n}{j+1} \binom{n}{j},$$

and

$$\tilde{Q}_a(n, j, 3) = \sum_{s \geq 1} \binom{n-2s-1}{j-3s-a} \binom{n+2s-1}{j+3s+a}.$$

*Proof.* First note that (13) is just a special case of Corollary 1 with  $k = 1$ .

The left hand side of (14) follows from Theorem 2. While the right hand side of (14) follows from Theorem 8, noting that

$$N(n, n-j) = N(n, j+1),$$

and

$$Q_a(n+1, n-j, 3) = \sum_{s \geq 1} \binom{n-2s-1}{n-j+s+a-1} \binom{n+2s-1}{n-j-s-a-1} = \tilde{Q}_a(n, j, 3),$$

using the symmetry of the Narayana numbers and the binomial coefficients.  $\square$

*Remark 2.* If  $j = 0, 1$ , then  $\tilde{Q}_a(n, j, 3) = 0$  for  $a = -1, 0, 1$ , and equation (14) follows immediately.

However, for  $j \geq 2$  there will be nonzero contributions from  $\tilde{Q}_a(n, j, 3)$  for  $a \leq j-3$ . For example, if  $j = 2$  then

$$\tilde{Q}_{-1}(n, 2, 3) = \sum_{s \geq 1} \binom{n-2s-1}{2-3s+1} \binom{n+2s-1}{2+3s-1} = \binom{n+1}{4},$$

and the right hand side of (14) becomes

$$\begin{aligned} & \frac{1}{n} \binom{n}{3} \binom{n}{2} - \binom{n+1}{4} = \frac{1}{n} \binom{n}{3} \binom{n}{2} - \binom{n}{4} - \binom{n}{3} \\ &= \binom{n}{3} \left[ \frac{n-1}{2} - 1 \right] - \binom{n}{4} = 2 \binom{n}{3} \left[ \frac{n-3}{4} \right] - \binom{n}{4} = \binom{n}{4} \end{aligned}$$

as expected.

More generally, we can use Theorem 12 when  $k \geq 2$ .

**Theorem 14.** For  $n \geq 1$ ,  $j \geq 0$ , and  $k \geq 2$  we have

$$\sum_{j \geq i_{k-1} \geq \dots \geq i_1 \geq 0} \left( \prod_{m=0}^{k-1} \binom{i_{m+2} + i_m}{2i_{m+1}} \right) \\ = N(n, j+1) - \left[ \tilde{Q}_1(n, j, k+2) - 2\tilde{Q}_0(n, j, k+2) + \tilde{Q}_{-1}(n, j, k+2) \right],$$

where

$$N(n, j+1) = \frac{1}{n} \binom{n}{j+1} \binom{n}{j},$$

and

$$\tilde{Q}_a(n, j, k+2) = \sum_{s \geq 1} \binom{n - (k+1)s - 1}{j - (k+2)s - a} \binom{n + (k+1)s - 1}{j + (k+2)s + a}.$$

*Proof.* From Corollary 1 we have

$$a_{n,231,j}^{(k)} = h_{k+2}(n+1, n-j).$$

The left hand side of Theorem 14 follows from Theorem 12. And the right hand side of Theorem 14 follows from Theorem 8, noting that

$$Q_a(n+1, n-j, k+2) = \sum_{s \geq 1} \binom{n - (k+1)s - 1}{n-j+s+a-1} \binom{n + (k+1)s - 1}{n-j-s-a-1} \\ = \tilde{Q}_a(n, j, k+2).$$

□

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